

COHERENT CHOICE FUNCTIONS, DESIRABILITY AND INDIFFERENCE

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ABSTRACT. We investigate how to model indifference with choice functions. We take the coherence axioms for choice functions proposed by Seidenfeld, Schervish and Kadane as a source of inspiration, but modify them to strengthen the connection with desirability. We discuss the properties of choice functions that are coherent under our modified set of axioms and the connection with desirability. Once this is in place, we present an axiomatisation of indifference in terms of desirability. On this we build our definition of indifference in terms of choice functions, which we discuss in some detail.

1. INTRODUCTION

The fundamental problem of study in decision theory is how to make a choice from within a set of available options. Rubín [21] proposed a theory of choice functions based on choices between more than two elements, where an agent determines which are the optimal ones for any input set. He started from von Neumann and Morgenstern's [33] axiomatisation of choices based on pairwise comparisons between the options.

One aspect of Rubín's axiomatisation is that the agent is assumed to be able to make a comparison between any pair of options (either one of them is strictly preferred to the other or they are indifferent). As a consequence, he is assumed to be indifferent between any two options chosen from a given set. However, the agent may be undecided between two options without being indifferent between them, particularly when the information available does not allow for a complete comparison of the options; this will for instance typically be the case when there is no relevant information available at all. This is one of the motivations for a theory of imprecise probabilities [34], where incomparability and indifference are distinguished. This point was also discussed by Kadane et al. [16] (and also [27]), who generalised Rubín's axioms [21] in order to allow for incomparability. One interesting aspect of their take on choice functions is that it does not require any predefined utility.

Imprecise probability is an umbrella term for mathematical models that are meant to be used in situations of imprecise or incomplete information, where it may not be possible (or advisable) to use (precise) probabilities. In particular, it covers sets of probability measures and various types of non-additive measures and functionals, such as coherent lower previsions, belief functions and possibility measures. All of these models can be expressed in terms of coherent sets of desirable gambles [19, 34, 35], which encode the gambles that a subject, whose beliefs we want to model, strictly prefers to the status quo. One of their advantages is that they avoid problems with conditioning on events of probability zero. They can be—and have been—used to replace probabilities in Bayesian networks, for predictive inference, and so on [5, 7, 11, 13, 18].

Sets of desirable gambles are typically a binary concept: they are characterised by pairwise comparisons between the available options, whereas in practice choice may be more complex. Choice functions based on binary comparisons only may represent behaviour that is considered to be paradoxical by some [3, 14, 22]. One of the aims of this paper is to

study how the more general—not necessarily binary—choice functions relate to the sets of desirable gambles that are now more commonly used in imprecise probabilities papers.

In order to establish such a relationship, we must pay some attention at this point to the options that subjects choose between on both approaches: *gambles* on the desirability approach, and *horse lotteries* in a choice function context. A detailed discussion of this issue can be found in a recent paper by Zaffalon and Miranda [36], but for our present purposes it suffices to mention that choosing between gambles essentially amounts to choosing between horse lotteries with only two rewards—say, winning a prize and not winning it—and the value of a gamble is essentially proportional to the probability of winning the prize on the corresponding horse lottery approach. This correspondence makes it relatively straightforward to translate the axioms proposed by Seidenfeld et al. [27] for choice functions on horse lotteries into their counterparts for gambles: we prefer to work with gambles here because that makes it easier to see the link with sets of desirable gambles. We will find it expedient to drop two particular axioms from the theory proposed by Seidenfeld et al. [27], thus making our coherence notion less restrictive. The first one is the Archimedean axiom, which is not fully compatible with the theory of coherent sets of desirable gambles—indeed it is letting go of Archimedean continuity that allows coherent sets of desirable gambles to avoid the problems that arise when conditioning on sets of probability zero. We have also not included the convexity axiom of Seidenfeld et al. [27], since it is incompatible with Walley–Sen maximality [30, 34], the optimality criterion typically associated with sets of desirable gambles. We discuss the axioms for our coherent choice functions and their main properties in Section 2.

In Section 3, we show that models of coherent sets of desirable gambles can be embedded within our theory of coherent choice functions, and we study the specific properties of choice functions that can be related to desirability. This exploratory work culminates in Section 4, where we focus on sets of probability measures and collections of maximal coherent sets of desirable gambles.

Once the connection between choice functions and desirability is in place, we move—in Section 5—to a definition and closer investigation of the property of indifference between two options. This is a crucial step towards enabling uncertainty models to represent symmetry [9]. For example, when modelling exchangeability and proving de Finetti-type representation theorems, an appropriate notion of indifference is essential. For choice functions, indifference was introduced by Seidenfeld [24], and for sets of desirable gambles by de Cooman and Quaeghebeur [11].

The connection between choice functions and sets of desirable gambles established in Section 3 leads us to investigate how to model choice that is compatible with a set of indifference assessments. Our approach to defining indifference focuses on the quotient space induced by the indifference relation. This allows us to exploit notions and techniques from linear algebra to deal with it mathematically. It also enables us to prove a representation result that makes it easy to consider conservative inference under indifference. We show how our notion is related to the one established by Seidenfeld [24].

We conclude the paper with some additional discussion in Section 6. In order not to impede the main flow of the argumentation, we have moved all proofs to an appendix.

2. CHOICE FUNCTIONS ON OPTION SETS

Consider a real vector space \mathcal{V} , provided with the vector addition $+$ and scalar multiplication. We denote by 0 the additive identity, or null vector. For any subsets O_1 and O_2 of \mathcal{V} and any λ in \mathbb{R} , we define $\lambda O_1 := \{\lambda u : u \in O_1\}$ and $O_1 + O_2 := \{u + v : u \in O_1, v \in O_2\}$.

Elements u of \mathcal{V} are intended as abstract representations of *options* amongst which a subject can express his preferences, by specifying, as we will see below, choice functions. Mostly, options will be real-valued maps on the possibility space, interpreted as uncertain rewards, and therefore also called *gambles*. Our reason for working with the more abstract notion of options—elements of some general vector space—is that in Section 5, we will need choice functions defined on *equivalence classes* of options. These equivalence classes again constitute a vector space—and hence are abstract options themselves—but they can no longer be interpreted easily and directly as gambles. Another reason for using options more general than gambles is that recent work [36] has shown that a very general theory of binary preference can be constructed using *vector-valued* gambles. These are no longer gambles in the above-mentioned sense, but still constitute a linear space.

We denote by $\mathcal{Q}(\mathcal{V})$ the set of all non-empty *finite* subsets of \mathcal{V} , a strict subset of the power set of \mathcal{V} . Elements O of $\mathcal{Q}(\mathcal{V})$ are the option sets amongst which a subject can choose his preferred options. When it is clear what vector space of options we are talking about, we will omit explicit mention of \mathcal{V} and simply write \mathcal{Q} .

Definition 1. A *choice function* C on \mathcal{Q} is a map

$$C: \mathcal{Q} \rightarrow \mathcal{Q} \cup \{\emptyset\}: O \mapsto C(O) \text{ such that } C(O) \subseteq O.$$

We collect all choice functions in the set \mathcal{C} .

The idea underlying this definition is that a choice function C selects the set $C(O)$ of ‘best’ options in the *option set* O , or, on another interpretation, the ones that cannot be rejected. Our definition resembles the one commonly used in the literature [1, 27, 29], except for a restriction to *finite* option sets, which, then again, is also not altogether unusual [15, 23, 28].

2.1. An equivalent representation. An equivalent notion of a choice function is that of a binary relation—called *choice relation*—on \mathcal{Q} .

Definition 2. A *choice relation* \triangleleft is a binary relation on \mathcal{Q} that satisfies the following two properties for all O, O_1 and O_2 in \mathcal{Q} :

- (i) if $O_2 \setminus O_1 \subseteq O \subseteq O_1 \cup O_2$ then $O_1 \triangleleft O_2 \Leftrightarrow O_1 \triangleleft O$;
- (ii) if $O_1 \cup O_2 \subseteq O$ then $(O_1 \triangleleft O \text{ and } O_2 \triangleleft O) \Leftrightarrow O_1 \cup O_2 \triangleleft O$.

We collect all choice relations in the set \mathcal{S} .

Given a choice function, we can use it to define a choice relation, and vice versa:

Definition 3. Given a choice function C , define the corresponding binary relation \triangleleft_C as

$$O_1 \triangleleft_C O_2 \Leftrightarrow C(O_1 \cup O_2) \subseteq O_2 \setminus O_1 \Leftrightarrow C(O_1 \cup O_2) \cap O_1 = \emptyset \text{ for all } O_1 \text{ and } O_2 \text{ in } \mathcal{Q}.$$

Conversely, given a choice relation \triangleleft , define the corresponding function C_\triangleleft as

$$C_\triangleleft(O) := \bigcap \{O' \subseteq O : O \setminus O' \triangleleft O\} = O \setminus \bigcup \{O' \subseteq O : O' \triangleleft O\} \text{ for all } O \text{ in } \mathcal{Q}.$$

The interpretation is clear: we find that $O_1 \triangleleft_C O_2$ if $u \notin C(O_1 \cup O_2)$ for all u in O_1 —in other words, when every option from O_1 is not chosen—rejected—when presented with $O_1 \cup O_2$ as option set. Furthermore, for C_\triangleleft , an option u is rejected from O —meaning that $u \in O$ and $u \notin C_\triangleleft(O)$ —if there is some $O' \subseteq O$ for which $u \in O'$ and $O' \triangleleft O$ —in other words, if there is some O' that contains u and such that O is preferred to O' .

Starting from a choice relation \triangleleft , it is clear from Definition 3 that C_\triangleleft is a choice function—meaning that $C_\triangleleft(O) \subseteq O$ for all O in \mathcal{Q} . We prove the converse argument—that \triangleleft_C satisfies the two properties of Definition 2—in the following proposition:

Proposition 1. *Given a choice function C , the corresponding binary relation \triangleleft_C is a choice relation.*

Interestingly, we can use C and \triangleleft interchangeably:

Proposition 2. *Given any choice function C , we have $C_{\triangleleft_C} = C$. Similarly, given any choice relation \triangleleft , we have $\triangleleft_{C_\triangleleft} = \triangleleft$. As a consequence, $\triangleleft_C = \triangleleft \Leftrightarrow C_\triangleleft = C$ for all C in \mathcal{C} and \triangleleft in \mathcal{S} .*

2.2. Rationality axioms. Seidenfeld et al. [27, Section 3] call a choice function C *coherent* if there is a non-empty set of probability-utility pairs \mathcal{S} such that $C(O)$ is the set of options in O that maximise expected utility for some element of \mathcal{S} . They also provide an axiomatisation for this type of coherence, based on the one for binary preferences [2]. One of their axioms is an ‘Archimedean’ continuity condition, and another one is a convexity condition, necessary for the connection with a set of probability-utility pairs.

We prefer to define coherence directly in terms of axioms, without reference to probabilities and utilities. In such a context, we see no compelling reason to adopt an Archimedean axiom, all the more so because we are interested in establishing the connection between choice functions and Walley’s [35] sets of desirable gambles, which need not satisfy this axiom. Furthermore, the convexity condition does not generally allow for choice functions that select the undominated options under some partial ordering, which is something we find natural, and will need later on.

For these reasons, we propose an alternative—and in some sense weaker—axiomatisation in Section 2.2.2 by dropping the Archimedean condition and the convexity condition of Seidenfeld et al. [27]. Then again, our second axiom is a slightly strengthened version of theirs, needed for the conditioning we intend to discuss in a later paper.

2.2.1. Some useful definitions. We call \mathbb{N} the set of all (strictly positive) integers, and $\mathbb{N}_0 := \mathbb{N} \cup \{0\}$. Also, we call $\mathbb{R}_{>0}$ the set of all (strictly) positive real numbers, and $\mathbb{R}_{\geq 0} := \mathbb{R}_{>0} \cup \{0\}$.

Given any subset O of \mathcal{V} , we define the *linear hull* $\text{span}(O)$ as the set of all finite linear combinations of elements of O :

$$\text{span}(O) := \left\{ \sum_{k=1}^n \lambda_k u_k : n \in \mathbb{N}, \lambda_k \in \mathbb{R}, u_k \in O \right\} \subseteq \mathcal{V},$$

the *positive hull* $\text{posi}(O)$ as the set of all positive finite linear combinations of elements of O :

$$\text{posi}(O) := \left\{ \sum_{k=1}^n \lambda_k u_k : n \in \mathbb{N}, \lambda_k \in \mathbb{R}_{>0}, u_k \in O \right\} \subseteq \text{span}(O),$$

and the *convex hull* $\text{CH}(O)$ as the set of convex combinations of elements of O :

$$\text{CH}(O) := \left\{ \sum_{k=1}^n \alpha_k u_k : n \in \mathbb{N}, \alpha_k \in \mathbb{R}_{\geq 0}, \sum_{k=1}^n \alpha_k = 1, u_k \in O \right\} \subseteq \text{posi}(O).$$

A subset O of \mathcal{V} is called a *convex cone* if it is closed under positive finite linear combinations, i.e. if $\text{posi}(O) = O$. A convex cone \mathcal{K} is called *proper* if $\mathcal{K} \cap -\mathcal{K} = \{0\}$.

With any proper convex cone $\mathcal{K} \subseteq \mathcal{V}$, we associate an ordering $\preceq_{\mathcal{K}}$ on \mathcal{V} , defined for all u and v in \mathcal{V} as follows:

$$u \preceq_{\mathcal{K}} v \Leftrightarrow v - u \in \mathcal{K} \Leftrightarrow 0 \preceq_{\mathcal{K}} v - u \Leftrightarrow u - v \preceq_{\mathcal{K}} 0.$$

We also write $u \succeq_{\mathcal{K}} v$ for $v \preceq_{\mathcal{K}} u$. The ordering $\preceq_{\mathcal{K}}$ is actually a *vector ordering*: it is a partial order—reflexive, antisymmetric and transitive—that satisfies the following two

characteristic properties:

$$u_1 \preceq_{\mathcal{K}} u_2 \Leftrightarrow u_1 + v \preceq_{\mathcal{K}} u_2 + v; \quad (1)$$

$$u_1 \preceq_{\mathcal{K}} u_2 \Leftrightarrow \lambda u_1 \preceq_{\mathcal{K}} \lambda u_2, \quad (2)$$

for all u_1, u_2, v in \mathcal{V} and λ in $\mathbb{R}_{>0}$. Conversely, given any vector ordering \preceq , the proper convex cone \mathcal{K} from which it is derived can always be retrieved by $\mathcal{K} = \{u \in \mathcal{V} : u \succeq 0\}$. When the abstract options are gambles, \preceq will typically be the pointwise order. When the options are equivalence classes, as in Section 5.2, the ordering will be typically the induced ordering on equivalence classes, defined in Equation (13) further on.

The vector space of options \mathcal{V} , ordered by the vector ordering $\preceq_{\mathcal{K}}$, is called an *ordered vector space* $(\mathcal{V}, \preceq_{\mathcal{K}})$. We will refrain from explicitly mentioning the actual proper convex cone \mathcal{K} we are using, and simply write \mathcal{V} to mean the ordered vector space, and use \preceq for the associated vector ordering.

Finally, with any vector ordering \preceq , we associate the strict partial ordering $<$ as follows:

$$u < v \Leftrightarrow (u \preceq v \text{ and } u \neq v) \Leftrightarrow v - u \in \mathcal{K} \setminus \{0\} \text{ for all } u, v \text{ in } \mathcal{V}.$$

We call u *positive* if $u > 0$, and collect all positive options in the convex cone $\mathcal{V}_{>0} := \mathcal{K} \setminus \{0\}$.

At this point, it is important to mention that from here on, we assume the cone \mathcal{K} to be fixed and given. It serves to identify a ‘positive orthant’ $\mathcal{V}_{>0} = \mathcal{K} \setminus \{0\}$, and should not be confused with the sets of desirable options (defined further on in Section 3.1), which are cones as well, and typically include the positive orthant $\mathcal{V}_{>0}$.

2.2.2. Rationality axioms for choice functions.

Definition 4. We call a choice function C on $\mathcal{Q}(\mathcal{V})$ *coherent* if for all O, O_1, O_2 in \mathcal{Q} , all u, v in \mathcal{V} and all λ in $\mathbb{R}_{>0}$:

- C₁. $C(O) \neq \emptyset$;
- C₂. if $u < v$ then $\{v\} = C(\{u, v\})$;
- C₃. a. if $C(O_2) \subseteq O_2 \setminus O_1$ and $O_1 \subseteq O_2 \subseteq O$ then $C(O) \subseteq O \setminus O_1$;
b. if $C(O_2) \subseteq O_1$ and $O \subseteq O_2 \setminus O_1$ then $C(O_2 \setminus O) \subseteq O_1$;
- C₄. a. if $O_1 \subseteq C(O_2)$ then $\lambda O_1 \subseteq C(\lambda O_2)$;
b. if $O_1 \subseteq C(O_2)$ then $O_1 + \{u\} \subseteq C(O_2 + \{u\})$.

We collect all coherent choice functions on \mathcal{V} in the set $\bar{\mathcal{C}}$.

Remark that Axiom C₃b is well defined, in the sense that none of C ’s arguments there can be empty. Indeed, $O_2 \setminus O$ is guaranteed to be non-empty: since $O \subseteq O_2 \setminus O_1$, we find that $O_2 \setminus O$ must include the set O_1 , which is non-empty by taking Axiom C₁ into account.

Parts C₃a and C₃b of Axiom C₃ are respectively known as *Sen’s condition α* and *Aizerman’s condition*. They are more commonly written as, respectively:

$$(O_1 \cap C(O_2) = \emptyset \text{ and } O_1 \subseteq O_2 \subseteq O) \Rightarrow O_1 \cap C(O) = \emptyset \text{ for all } O, O_1, O_2 \text{ in } \mathcal{Q}, \quad (3)$$

and

$$(O_1 \cap C(O_2) = \emptyset \text{ and } O \subseteq O_1) \Rightarrow O_1 \cap C(O_2 \setminus O) = \emptyset \text{ for all } O, O_1, O_2 \text{ in } \mathcal{Q}. \quad (4)$$

The first one means that any option that is rejected in a option set O_2 must also be rejected in any larger option set O (i.e., the inclusion of new options in the set cannot make an option that was rejected to become acceptable), while the second one means that if we reject an option in O_2 we should also reject it when we remove some options that are also rejected from the option set (i.e., the removal of rejected options does not alter the judgments about the remaining ones).

The following proposition collects a number of immediate properties of coherent choice functions.

Proposition 3. *The following statements hold for any coherent choice function C :*

- (i) $\lambda C(O) + \{u\} = C(\lambda O + \{u\})$ for all O in \mathcal{Q} , λ in $\mathbb{R}_{>0}$ and u in \mathcal{V} ;
- (ii) for all u_1, u_2 in \mathcal{V} such that $u_1 \leq u_2$, all O in \mathcal{Q} and all v in $O \setminus \{u_1, u_2\}$:
 - a. if $u_2 \in O$ and $v \notin C(O \cup \{u_1\})$ then $v \notin C(O)$;
 - b. if $u_1 \in O$ and $v \notin C(O)$ then $v \notin C(\{u_2\} \cup O \setminus \{u_1\})$;
- (iii) C is insensitive to the omission of non-chosen options [12, Definition 11]: $C(O') = C(O)$ for all O, O' in \mathcal{Q} such that $C(O) \subseteq O' \subseteq O$;
- (iv) $C(C(O)) = C(O)$ for all O in \mathcal{Q} .

As already announced and motivated earlier, our axioms in Definition 4 are a subset of the ones considered by Seidenfeld et al. [27], who use *horse lotteries* as the basic options a subject must choose between. The set of all horse lotteries plays a similar role there as our vector space of options \mathcal{V} does here. We omit two of Seidenfeld et al.'s [27] axioms: the Archimedean axiom, ensuring some continuity of coherent choice functions, and a convexity axiom. However, in further work we do intend to study the implications of adding the convexity axiom in the context of the relation between choice functions and sets of desirable options.

Our Axiom C_2 is a slightly strengthened version of one of Seidenfeld et al.'s [27] axioms. To compensate for this slightly weaker requirement, they impose the two properties (ii)a and (ii)b of our Proposition 3 as separate rationality axioms [27, Axiom 4]. Our proofs for them rely quite heavily on, amongst other things, our slightly stronger Axiom C_2 . Our reason for using this strengthened version is that we want to concentrate on the relationship with desirability: any choice function that does not satisfy Axiom C_2 does not induce a coherent set of desirable gambles, as we will see in Section 3.3.

Seidenfeld et al.'s [27] call a choice function C , defined on the set $\mathcal{Q}(\mathcal{H})$ of all finite subsets of the set \mathcal{H} of all horse lotteries, coherent if

- C_1^* . $C(O) \neq \emptyset$ for all O in $\mathcal{Q}(\mathcal{H})$;
- C_2^* . for all O in $\mathcal{Q}(\mathcal{H})$ and H, H_1 and H_2 in \mathcal{H} such that $H \notin \{H_1, H_2\}$ and $H_1 \leq H_2$:
 - a. if $H_2 \in O$ and $H \notin C(\{H_1\} \cup O)$ then $H \notin C(O)$;
 - b. if $H_1 \in O$ and $H \notin C(O)$ then $H \notin C(\{H_2\} \cup O \setminus \{H_1\})$;
- C_3^* . for all O, O_1 and O_2 in $\mathcal{Q}(\mathcal{H})$:
 - a. if $C(O_2) \subseteq O_2 \setminus O_1$ and $O_1 \subseteq O_2 \subseteq O$ then $C(O) \subseteq O \setminus O_1$;
 - b. if $C(O_2) \subseteq O_1$ and $O \subseteq O_2 \setminus O_1$ then $C(O_2 \setminus O) \subseteq O_1$;
- C_4^* . $O_1 \triangleleft_C O_2 \Leftrightarrow (\alpha O_1 + (1 - \alpha)\{H\}) \triangleleft_C (\alpha O_2 + (1 - \alpha)\{H\})$, for all O_1 and O_2 in $\mathcal{Q}(\mathcal{H})$ and H in \mathcal{H} ;
- C_5^* . if $O \subseteq O_1 \subseteq CH(O)$ then $C(O) \subseteq C(O_1)$, for all O and O_1 in $\mathcal{Q}(\mathcal{H})$;
- C_6^* . for all O, O', O'' in $\mathcal{Q}(\mathcal{H})$ and all sequences O_i' and O_i'' in $\mathcal{Q}(\mathcal{H})$ such that O_i' converges point-wise to O' and O_i'' converges point-wise to O'' (for i in \mathbb{N}):
 - a. if $(\forall i \in \mathbb{N}) O_i'' \triangleleft_C O_i'$ and $O' \triangleleft_C O$ then $O'' \triangleleft_C O$;
 - b. if $(\forall i \in \mathbb{N}) O_i'' \triangleleft_C O_i'$ and $O \triangleleft_C O''$ then $O \triangleleft_C O''$;

where \leq is some predetermined reflexive order on \mathcal{H} .

Seidenfeld et al.'s [27] choice functions are defined on horse lotteries while we define our choice functions on any linear space, so direct comparison is not immediate. But there is a very close formal resemblance between the two sets of axioms. Indeed, Axioms C_1 and C_1^* correspond; so do Axioms C_3 and C_3^* ; and—perhaps less obviously—Axioms C_4 and C_4^* . Axiom C_2^* , being a dominance condition, is essentially Proposition 3(ii), and can be regarded as a weaker counterpart of Axiom C_2 . In the present setting, we are not interested

in Axiom C_5^* because it is in general not compatible with maximality as a decision rule; neither are we interested in Axiom C_6^* because it is a continuity requirement that is not fully compatible with sets of desirable gambles.

2.2.3. *Rationality axioms for choice relations.* It turns out we can formulate the rationality axioms for choice functions—Definition 4—directly in terms of choice relations.

Definition 5. We call a choice relation \triangleleft on $\mathcal{Q}(\mathcal{V})$ *coherent* if for all O, O_1, O_2 in \mathcal{Q} , all u in \mathcal{V} and all λ in $\mathbb{R}_{>0}$:

- S₁. $O \not\triangleleft O$;
- S₂. if $u < v$ then $\{u\} \triangleleft \{v\}$;
- S₃. a. if $O_1 \triangleleft O_2$ and $O_2 \subseteq O$ then $O_1 \triangleleft O$;
b. if $O_1 \triangleleft O_2$ and $O \subset O_1$ then $O_1 \setminus O \triangleleft O_2 \setminus O$;
- S₄. a. if $O_1 \triangleleft O_2$ then $\lambda O_1 \triangleleft \lambda O_2$;
b. if $O_1 \triangleleft O_2$ then $O_1 + \{u\} \triangleleft O_2 + \{u\}$.

We collect all coherent choice relations on \mathcal{V} in the set $\bar{\mathcal{S}}$.

The following proposition collects a number of immediate properties of coherent choice relations.¹

Proposition 4. *The following statements hold for any coherent choice relation \triangleleft :*

- (i) \triangleleft is a strict partial order: it is irreflexive and transitive;
- (ii) for all O_1 and O_2 in \mathcal{Q} , all u in \mathcal{V} and λ in $\mathbb{R}_{>0}$: $O_1 \triangleleft O_2 \Leftrightarrow \lambda O_1 + \{u\} \triangleleft \lambda O_2 + \{u\}$.

The rationality axioms of Definitions 4 and 5 correspond:

Proposition 5. *Given any coherent choice function C , \triangleleft_C is a coherent choice relation. Conversely, given any coherent choice relation \triangleleft , C_\triangleleft is a coherent choice function.*

2.3. **The ‘is not more informative than’ relation.** Because we are interested in *conservative reasoning* with choice functions, we want to look for those implications of a given assessment that are as ‘uninformative’ as possible. To this end, we need to introduce some binary relation \sqsubseteq on the set \mathcal{C} of all choice functions, having the specific interpretation of ‘not more informative than’, or, in other words, ‘at least as uninformative as’.

Definition 6. Given two choice functions C_1 and C_2 in \mathcal{C} , we call C_1 *not more informative than* C_2 —and we write $C_1 \sqsubseteq C_2$ —if $(\forall O \in \mathcal{Q}) C_1(O) \supseteq C_2(O)$.

This intuitive way of ordering choice functions is also used by Bradley [4] and Van Camp et al. [31]. The underlying idea is that a choice function is more informative when it consistently chooses more specifically—or more restrictively—amongst the available options. The corresponding notion for choice relations is as follows: given two choice relations \triangleleft_1 and \triangleleft_2 , \triangleleft_1 is not more informative than \triangleleft_2 —we write $\triangleleft_1 \sqsubseteq_{\mathcal{S}} \triangleleft_2$ —if $\triangleleft_1 \subseteq \triangleleft_2$, or in other words, if $O_1 \triangleleft_1 O_2$ implies $O_1 \triangleleft_2 O_2$, for all O_1 and O_2 in \mathcal{Q} . That is, choice relation \triangleleft_2 is at least as informative as \triangleleft_1 when every comparison made by \triangleleft_1 is also made by \triangleleft_2 .

Since, by definition, \sqsubseteq is a product ordering of set inclusions, the following result is immediate [6].

Proposition 6. *The structure $(\mathcal{C}; \sqsubseteq)$ is a complete lattice:*

- (i) *it is a partially ordered set, or poset, meaning that the binary relation \sqsubseteq on \mathcal{C} is reflexive, antisymmetric and transitive;*

¹See [16] for a similar result.

- (ii) for any subset \mathcal{C}' of \mathcal{C} , its infimum $\inf \mathcal{C}'$ and its supremum $\sup \mathcal{C}'$ with respect to the ordering \sqsubseteq exist in \mathcal{C} , and are given by $\inf \mathcal{C}'(O) = \bigcup_{C \in \mathcal{C}'} C(O)$ and $\sup \mathcal{C}'(O) = \bigcap_{C \in \mathcal{C}'} C(O)$ for all O in \mathcal{Q} .

Moreover, $(\mathcal{S}; \sqsubseteq_{\mathcal{S}})$ is a complete lattice as well.

The idea underlying these notions of infimum and supremum is that $\inf \mathcal{C}'$ is the most informative model that is not more informative than any of the models in \mathcal{C}' , and $\sup \mathcal{C}'$ the least informative model that is not less informative than any of the models in \mathcal{C}' .

We will also consider the poset $(\bar{\mathcal{C}}; \sqsubseteq)$ of all coherent choice functions, where $\bar{\mathcal{C}} \sqsubseteq \mathcal{C}$ inherits the partial order \sqsubseteq from \mathcal{C} . Similarly, in the poset $(\bar{\mathcal{S}}; \sqsubseteq_{\mathcal{S}})$ of all coherent choice relations, $\bar{\mathcal{S}}$ is assumed to inherit the partial order $\sqsubseteq_{\mathcal{S}}$ from \mathcal{S} .

Also from an order-theoretic point of view, choice functions or choice relations are equivalent:

Proposition 7. *The posets $(\mathcal{C}; \sqsubseteq)$ and $(\mathcal{S}; \sqsubseteq_{\mathcal{S}})$ are order isomorphic [6]: there is a map f from \mathcal{C} to \mathcal{S} such that $C_1 \sqsubseteq C_2 \Leftrightarrow f(C_1) \sqsubseteq_{\mathcal{S}} f(C_2)$ for all C_1 and C_2 in \mathcal{C} . Moreover, the posets $(\bar{\mathcal{C}}; \sqsubseteq)$ and $(\bar{\mathcal{S}}; \sqsubseteq_{\mathcal{S}})$ are order isomorphic as well.*

This proposition implies that we can regard \sqsubseteq and $\sqsubseteq_{\mathcal{S}}$ essentially as the same partial orders, each defined on their respective domains, in the sense that each of the partial orders can be obtained from the other one just by renaming. Therefore, from now on, we will identify $\sqsubseteq_{\mathcal{S}}$ with \sqsubseteq , and only use the latter.

Proposition 8. *$(\bar{\mathcal{C}}; \sqsubseteq)$ is a complete infimum-semilattice: $\bar{\mathcal{C}}$ is closed under arbitrary non-empty infima, so $\inf \mathcal{C}' \in \bar{\mathcal{C}}$ for any non-empty subset \mathcal{C}' of $\bar{\mathcal{C}}$.*

As a result of Proposition 7, $(\bar{\mathcal{S}}; \sqsubseteq)$ is a complete infimum-semilattice as well.

This, for instance, allows us to do conservative inference with choice functions: if we consider that there is a coherent choice function that represents a subject's beliefs (coherent choices) and we can only tell that it belongs to a family $\{C_i : i \in I\}$, the conservative option is to consider its infimum $\inf_{i \in I} C_i$. This choice function is still guaranteed to be coherent, and as a consequence, it satisfies all the rationality requirements discussed above.

As we will show in Proposition 14 further on, there is a least informative coherent choice function, namely the infimum of all coherent choice functions: it is given by $C(O) = \{u \in O : (\forall v \in O) u \not\prec v\}$ for all O in \mathcal{Q} .

3. THE RELATION WITH SETS OF DESIRABLE OPTIONS

In general, a choice function cannot be characterised using only pairwise comparisons of options, meaning that a binary relation on options does not uniquely determine a choice function. In this section, we study a special class of choice functions that *are* determined by their restrictions to option sets of cardinality two.

3.1. Sets of desirable options. Sets of desirable options are a(n obvious and immediate) generalisation of *sets of desirable gambles*. Gambles are real-valued maps on a possibility space \mathcal{X} , interpreted as uncertain rewards. They can be seen as vectors in the vector space $\mathbb{R}^{\mathcal{X}}$. Here we generalise this notion by looking at a general (abstract) vector space \mathcal{V} of (abstract) options, rather than gambles, provided with a vector ordering \leq ; see also Section 2.2.1. We will see that sets of desirable options amount to a pairwise comparison of options and therefore correspond to a special kind of choice functions.

A set of desirable options D is simply a subset of the vector space of options \mathcal{V} —its interpretation will be that it contains those options that some subject strictly prefers to the

status quo 0. We collect all possible such sets of desirable options in the set \mathcal{D} . As we did for choice functions, we pay special attention to *coherent* sets of desirable options. The following is an immediate generalisation of existing coherence definitions [7, 10, 11, 19] from gambles to abstract options.

Definition 7. A set of desirable options D is called coherent if for all u and v in \mathcal{V} and λ in $\mathbb{R}_{>0}$:

D₁. $0 \notin D$;

D₂. $\mathcal{V}_{>0} \subseteq D$;

D₃. if $u \in D$ then $\lambda u \in D$;

D₄. if $u, v \in D$ then $u + v \in D$.

We let $\bar{\mathcal{D}}$ denote the set of all coherent sets of desirable options.

Axioms D₃ and D₄ turn coherent sets of desirable options D into convex cones—meaning that $\text{posi}(D) = D$. They include the positive options due to Axiom D₂, but not the zero option due to Axiom D₁. As an immediate consequence, their intersection with $\mathcal{V}_{<0} := -\mathcal{V}_{>0}$ is empty. As usual, we may associate with the convex cone D a strict partial order \prec —called *preference relation*—on \mathcal{V} , by letting $u \prec v \Leftrightarrow 0 \prec v - u \Leftrightarrow v - u \in D$, leading to $D = \{u \in \mathcal{V} : 0 \prec u\}$ [11, 19]. The connection between this order and the one associated to choice functions will be clarified in Section 3.3.

3.2. The ‘is not more informative than’ relation. Sets of desirable options can be ordered according to a ‘not more informative than’ relation, analogously to the ordering introduced for choice functions.

Definition 8. Given two sets of desirable options D_1, D_2 in \mathcal{D} , we call D_1 *not more informative than* D_2 when $D_1 \subseteq D_2$.

Because the ordering of sets of desirable options \subseteq is just set inclusion, it is a partial ordering on \mathcal{D} , and the poset $(\mathcal{D}; \subseteq)$ is a complete lattice, with supremum operator \cup , and infimum operator \cap . Next we investigate the structure of the set of all coherent sets of desirable options:

Proposition 9. $(\bar{\mathcal{D}}; \subseteq)$ is a complete infimum-semilattice, or alternatively, $\bar{\mathcal{D}}$ is an intersection structure—closed under arbitrary non-empty intersections. The least informative (smallest) set of desirable options D_v is given by $D_v := \mathcal{V}_{>0}$.

We will refer to D_v as the *vacuous set of desirable options*.

It will be useful to also consider the maximally informative, or *maximal*, coherent sets of desirable options.² They are the undominated elements of the complete infimum-semilattice $(\bar{\mathcal{D}}; \subseteq)$; we collect them into a set $\hat{\mathcal{D}}$:

$$\hat{\mathcal{D}} := \{D \in \bar{\mathcal{D}} : (\forall D' \in \bar{\mathcal{D}})(D \subseteq D' \Rightarrow D = D')\}.$$

First, we prove a useful proposition that will allow us to characterise these maximal elements very elegantly.

Proposition 10. *Given any coherent set of desirable options D and any non-zero option $u \notin D$, $\text{posi}(D \cup \{-u\})$ is a coherent set of desirable options. As a consequence, a coherent set of desirable options D is maximal if and only if*

$$(\forall u \in \mathcal{V} \setminus \{0\})(u \in D \text{ or } -u \in D). \quad (5)$$

²The discussion in the rest of this section is based on similar discussions about sets of desirable gambles [5, 11, 20]. We repeat the details here *mutatis mutandis* to make the paper more self-contained.

Next we show that the set of all coherent sets of desirable options is dually atomic, meaning that any coherent set of desirable options is the infimum of its non-empty set of dominating maximal coherent sets of desirable options:

Proposition 11. *For any coherent set of desirable options D , its set of dominating maximal coherent sets of desirable options $\hat{D}_D := \{\hat{D} \in \hat{\mathcal{D}} : D \subseteq \hat{D}\}$ is non-empty. Moreover, $D = \inf \hat{D}_D$.*

3.3. Connection between choice functions and sets of desirable options. We now set out to establish a connection between choice functions and sets of desirable options.

Definition 9. Given a choice function C , we say that an option v is *chosen over* some option u whenever $u \notin C(\{u, v\})$, or equivalently, whenever $\{u\} \triangleleft_C \{v\}$. Similarly, given a set of desirable options D , we say that an option v is *preferred to* some option u whenever $v - u \in D$, or equivalently, $u \triangleleft v$. We call a choice function C and a set of desirable options D *compatible* when

$$u \notin C(\{u, v\}) \Leftrightarrow v - u \in D \Leftrightarrow u \triangleleft v \text{ for all } u, v \in \mathcal{V}.$$

Compatibility means that the behaviour of the choice function *restricted to pairs of options* reflects the behaviour of the set of desirable options.³ So, a choice function C will have at most one compatible set of desirable options, whereas conversely, a set of desirable options D may have many compatible choice functions: compatibility only directly influences the behaviour of a choice function on pairs.

The definition of compatibility with desirability is clearer in terms of the relations: a choice relation \triangleleft is compatible with some preference relation \triangleleft if $\{u\} \triangleleft \{v\} \Leftrightarrow u \triangleleft v$ for all u and v in \mathcal{Q} . Loosely speaking, \triangleleft and \triangleleft are compatible if \triangleleft lifts \triangleleft from options to sets of options. This observation also sheds light on the necessity of our Axiom C_2 , which Seidenfeld et al. [27] use a weakened version of. Indeed, if a choice relation does not satisfy Axiom C_2 then necessarily $0 \not\triangleleft u$ for some $0 < u$, contradicting Axiom D_2 .

3.3.1. From choice functions to desirability. We begin by studying the properties of the set of desirable options compatible with a given coherent choice function. Since compatibility is related to the restriction of the choice function to pairwise comparison, it is not surprising that each choice function has a unique compatible set of desirable options:

Proposition 12. *For any coherent choice function C in $\bar{\mathcal{C}}$, the unique compatible coherent set of desirable options D_C is given by $D_C := \{u \in \mathcal{V} : 0 \notin C(\{0, u\})\} = \{u \in \mathcal{V} : \{0\} \triangleleft_C \{u\}\}$.*

Example 1. Consider $\mathcal{V} = \mathbb{R}^2$, and let us denote the two components of any vector u in \mathbb{R}^2 by u_1 and u_2 , so $u = (u_1, u_2)$. Let C be the choice function given by

$$C(O) := \{u \in O : \text{there is no } v \in O \text{ such that } v_1 + v_2 > u_1 + u_2\}. \quad (6)$$

It is easy to check that C is a coherent choice function—it is an example of what we will call an E -admissible choice function later. And

$$D_C = \{u \in \mathcal{V} : u_1 + u_2 > 0\} \quad (7)$$

is its associated coherent set of desirable options. \square

³See Seidenfeld et al. [25] for an axiomatisation of imprecise preferences in the context of binary comparisons of horse lotteries, rather than gambles.

3.3.2. *From desirability to choice functions.* We collect in \bar{C}_D all the coherent choice functions that are compatible with a given coherent set of desirable options D :

$$\begin{aligned}\bar{C}_D &:= \{C \in \bar{C} : (\forall u, v \in \mathcal{V})(v \notin C(\{u, v\}) \Leftrightarrow u - v \in D)\} \\ &= \{C \in \bar{C} : D_C = D\}.\end{aligned}$$

The correspondence between choice functions and sets of desirable options can be many-to-one, in the sense that several different coherent choice functions may be compatible with the same coherent set of desirable options. An important role is played by the most uninformative of them:

Proposition 13. *Given a coherent set of desirable options D , the infimum—most uninformative element— $\inf \bar{C}_D$ of its set of compatible coherent choice functions \bar{C}_D is the coherent choice function C_D , defined by*

$$\begin{aligned}C_D(O) &:= \{u \in O : (\forall v \in O)v - u \notin D\} \\ &= \{u \in O : (\forall v \in O)u \not\prec v\} \text{ for all } O \text{ in } \mathcal{Q}.\end{aligned}\tag{8}$$

The coherent choice function C_D is the least informative choice function that is compatible with a coherent set of desirable options D : it is based on the binary ordering represented by D and nothing else. As we will see in Proposition 19, there typically are other coherent choice functions C compatible with D , but they encode more information than just the binary ordering represented by D . Proposition 13 is especially interesting because it shows that the most conservative choice function based on a strict partial order of options, is the choice function based on *maximality*⁴—the one that selects the *undominated* options under the strict partial order \prec associated with a coherent set of desirable options D . Any choice function that is based on maximality under such a strict partial order is coherent.

Proposition 13 can also be interpreted in terms of choice relations: given a coherent set of desirable gambles D —or a coherent preference relation \prec —, the most uninformative choice relation \triangleleft_D that is compatible with it, is by Proposition 7 equal to \triangleleft_{C_D} , and determined by

$$O_1 \triangleleft_D O_2 \Leftrightarrow (\forall u \in O_1)(\exists v \in O_2)u \prec v, \text{ for all } O_1 \text{ and } O_2 \text{ in } \mathcal{Q}.$$

This shows that the partial order \triangleleft_D is simply the result of lifting \prec from elements to sets.

Example 2. Let D be the coherent set of desirable options given by Equation (7). Its associated coherent choice function C_D is given by

$$C(O) := \{u \in O : \text{there is no } v \in O \text{ such that } v_1 + v_2 > u_1 + u_2\},$$

so it coincides with the choice function given by Equation (6). \square

Although in this particular example the two procedures coincide, there is in general more than one coherent choice function that is compatible with a coherent set of desirable options. In other words, a coherent choice function is not uniquely determined by its restriction to binary comparisons; we will give an example in Example 6 further on.

Proposition 8 guarantees that there is a unique smallest—least informative—coherent choice function. We will call it the *vacuous choice function*, and denote it by C_v .

Proposition 14. *The vacuous choice function C_v is given by $C_v(O) = C_{D_v}(O) = \{u \in O : (\forall v \in O)u \not\prec v\}$ for all O in \mathcal{Q} . It selects from any set of options the ones that are undominated under the strict vector ordering \prec .*

⁴Note that maximality here refers to the optimality decision criterion [30, 34], and it does not mean that the set of desirable options D is a maximal one, in the sense considered in Section 3.2.

Similarly, the vacuous choice relation \triangleleft_v is equal to \triangleleft_{C_v} by Proposition 7, and is determined by $O_1 \triangleleft_v O_2 \Leftrightarrow (\forall u \in O_1)(\exists v \in O_2)u < v$ for all O_1 and O_2 in \mathcal{Q} .

Example 3. Consider, as a simple example, the case that the vector ordering is total, meaning that for any u, v in \mathcal{V} , either $u < v$, $v < u$ or $u = v$. It then follows from Proposition 14 that, for any coherent choice function C , $C(O) \subseteq C_v(O) = \max O$ for all $O \in \mathcal{Q}$, where $\max O$ is the unique largest element of the finite option set O according to the strict total ordering $<$. But then Axiom C_1 guarantees that $C(O) = C_v(O) = \max O$ for all $O \in \mathcal{Q}$, so C_v is the *only* coherent choice function. \square

3.3.3. Properties of the relation between choice functions and desirability. Since sets of desirable options represent only pairwise comparison, and are therefore generally less expressive than choice functions, we expect that going from a choice function to a compatible set of desirable options leads to a loss of information, whereas going the opposite route does not. This is confirmed by Propositions 15 and 16, and in particular by their Corollary 17.

Proposition 15. Consider any set of coherent choice functions $\mathcal{C}' \subseteq \bar{\mathcal{C}}$. Then $D_{\inf \mathcal{C}'} = \inf\{D_C : C \in \mathcal{C}'\}$ and $C_{\inf\{D_C : C \in \mathcal{C}'\}} \subseteq \inf \mathcal{C}'$, and therefore also $C_{D_{\inf \mathcal{C}'}} \subseteq \inf \mathcal{C}'$.

Definition 10. For any set of coherent sets of desirable options $\mathcal{D}' \subseteq \bar{\mathcal{D}}$, we define the ‘infimum of maximality’ choice function as $C_{\mathcal{D}'} := \inf\{C_D : D \in \mathcal{D}'\}$.

Proposition 16. Consider any set of coherent sets of desirable options $\mathcal{D}' \subseteq \bar{\mathcal{D}}$. Then $D_{C_{\mathcal{D}'}} = \inf \mathcal{D}'$. Moreover, $C_{\inf \mathcal{D}'} \subseteq C_{\mathcal{D}'}$.

This result also helps to highlight that coherent choice functions are indeed more informative than coherent sets of desirable options, in the sense that two different coherent choice functions ($C_{\inf \mathcal{D}'}$ and $C_{\mathcal{D}'}$) may determine the same coherent set of desirable options ($\inf \mathcal{D}'$) when restricted to option sets of cardinality two. Thus, we need to move to this richer model in order to fully represent the beliefs encompassed by $\{C_D : D \in \mathcal{D}'\}$. Moreover, this happens even if each of the choice functions in the set \mathcal{D}' is fully determined by a coherent set of desirable options.

From these two results we immediately deduce the following:

Corollary 17. Consider any coherent set of desirable options $D \in \bar{\mathcal{D}}$ and any coherent choice function $C \in \bar{\mathcal{C}}$. Then $D = D_{C_D}$ and $C_{D_C} \subseteq C$.

If we consider the maps

$$D_\bullet : \mathcal{D} \rightarrow \mathcal{C} : D \mapsto C_D \text{ as defined by Proposition 12}$$

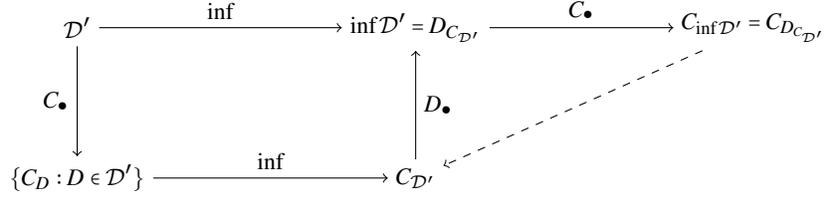
$$C_\bullet : \mathcal{C} \rightarrow \mathcal{D} : C \mapsto D_C \text{ as defined in Equation (8),}$$

then Corollary 17 essentially states that $D_\bullet \circ C_\bullet = \text{id}$, the identity map, while $(C_\bullet \circ D_\bullet)(C) \subseteq C$ for all C in $\bar{\mathcal{C}}$. This observation helps us interpret Propositions 15 and 16 using the commuting diagrams for Proposition 15:

$$\begin{array}{ccc} \mathcal{C}' & \xrightarrow{\text{inf}} & \inf \mathcal{C}' \\ \downarrow D_\bullet & & \downarrow D_\bullet \\ \{D_C : C \in \mathcal{C}'\} & \xrightarrow{\text{inf}} & \inf\{D_C : C \in \mathcal{C}'\} = D_{\inf \mathcal{C}'} \xrightarrow{C_\bullet} C_{\inf\{D_C : C \in \mathcal{C}'\}} = C_{D_{\inf \mathcal{C}'}} \end{array}$$

← (dashed arrow from $\inf \mathcal{C}'$ to $C_{\inf\{D_C : C \in \mathcal{C}'\}}$)

and for Proposition 16:



Full lines indicate the maps (D_\bullet , C_\bullet or inf), while dashed lines indicate an “is not more informative than” (Ξ for choice functions, \subseteq for sets of desirable gambles) relation.

Example 5 in Section 4 further on will show that the inequalities in these results can be strict; this does not seem surprising for the inequality of Proposition 15, but is perhaps less intuitively obvious for the one in Proposition 16.

4. OTHER TYPES OF COHERENT CHOICE FUNCTIONS

Given a coherent set of desirable options, we can define a coherent choice function by selecting undominated elements as in Equation (8). However, these are not the only possible coherent choice functions: for instance, any infimum of such coherent choice functions is still coherent.

Proposition 18. *Consider any set of coherent sets of desirable options $\mathcal{D}' \subseteq \tilde{\mathcal{D}}$, then $C_{\mathcal{D}'}$ is a coherent choice function.*

This proposition gives a procedure to define coherent choice functions complying with a sensitivity analysis interpretation, where the underlying uncertainty model is that of coherent sets of desirable options: we may consider a set \mathcal{D}' of possible models and then the set of coherent choice functions they determine; by taking the infimum of this set we end up choosing those options that are considered acceptable by at least one of the possible models.

Example 4. Consider again $\mathcal{V} = \mathbb{R}^2$, and let D_1, D_2 the coherent sets of desirable options given by

$$D_1 = \{u \in \mathbb{R}^2 : u(x) + u(y) > 0\} \text{ and } D_2 = \{u \in \mathbb{R}^2 : 2u(x) + u(y) > 0\}.$$

Then the choice function $C = \text{inf}\{C_{D_1}, C_{D_2}\}$ is coherent. If we consider for instance the option set $O = \{(4, -4), (2, -1), (0, 0), (-1, 2)\}$, we obtain

$$C(O) = \text{inf}\{C_{D_1}(O), C_{D_2}(O)\} = \{(2, -1), (-1, 2), (4, -4)\}.$$

The first two options are elements of $C_{D_1}(O)$ and the last one is an element of $C_{D_2}(O)$. \square

As special cases, we may consider ‘infimum of maximality’ choice functions where some additional condition is imposed on the coherent sets of desirable options. We will investigate two such situations in this section, although other are possible. In Definition 11, we focus only on sets of *maximal* coherent sets of desirable options, that we introduced in Section 3.2. Their interest lies in the fact that they are related to mass functions, as we discussed quite thoroughly in [5, 11, 20] in the context of gambles.

Definition 11. If $\mathcal{D}' \subseteq \hat{\mathcal{D}}$ is a set of *maximal* coherent set of desirable options, the coherent choice function $C_{\mathcal{D}'}$ is called *M-admissible*. We will also denote it by $C_{\mathcal{D}'}^M$, as a reminder that the infimum is taken over maximal sets.

In particular, we can consider the M-admissible choice functions for the set $\mathcal{D}' = \hat{\mathcal{D}}$ of all maximal coherent set of desirable options that include a coherent set of desirable options D .

In order not to overburden the notation, we let

$$C_D^M := C_{\hat{\mathcal{D}}_D}^M = \inf\{C_{\hat{D}} : \hat{D} \in \hat{\mathcal{D}} \text{ and } D \subseteq \hat{D}\}, \quad (9)$$

and similarly to what we did before, we introduce the map

$$C_{\bullet}^M : \mathcal{D}' \rightarrow \mathcal{C}' : D \mapsto C_D^M \text{ as defined in Equation (9).}$$

The following result can be regarded as a particular case of Proposition 16, where all the coherent sets of desirable options are maximal ones. As we have seen there too, the diagram below commutes if we focus on sets of desirable options, but this is not longer the case if we consider the more informative model of coherent choice functions.

$$\begin{array}{ccc} & D' & \\ & \uparrow & \searrow C_{\bullet} \\ C_{\bullet}^M & & \\ & \downarrow & \\ & C_{D'}^M & \leftarrow \text{---} C_{D'} \end{array}$$

Proposition 19. *Consider any coherent set of desirable options $D' \in \bar{\mathcal{D}}$. Then $D' = D_{C_{D'}^M}$ and $C_{D'} \subseteq C_{D'}^M$.*

The inequality in Proposition 19 can be strict—meaning that $C_{D'} \subset C_{D'}^M$ for some coherent set of desirable options D' —as we will show in Example 6 below.

As another special case, we consider choice functions associated with Levi's notion of E-admissibility [17, Chapter 5]. They are based on a non-empty *set of mass functions*. Consider a *finite* possibility space \mathcal{X} , and maps from \mathcal{X} to \mathbb{R} (the options are *gambles* now), forming the vector space $\mathcal{V} = \mathbb{R}^{\mathcal{X}}$ of *finite* dimension $|\mathcal{X}|$. The vector ordering \leq we associate with this vector space of options is the pointwise ordering of real maps: $u \leq v \Leftrightarrow (\forall x \in \mathcal{X}) u_x \leq v_x$, where, for instance, $u_x = u(x)$ is the x -component of the option u . We call any map $p: \mathcal{X} \rightarrow \mathbb{R}$ with $(\forall x \in \mathcal{X}) p(x) \geq 0$ and $\sum_{x \in \mathcal{X}} p(x) = 1$ a (probability) *mass function*, and we associate an *expectation* E_p with p by letting $E_p(u) := \sum_{x \in \mathcal{X}} p(x) u_x$ for all u in \mathcal{V} . We collect in $\Sigma_{\mathcal{X}}$ all the mass functions on \mathcal{X} .

With a mass function p , we associate a set of desirable options

$$D_p := \mathcal{V}_{>0} \cup \{u \in \mathcal{V} : E_p(u) > 0\} \quad (10)$$

and a choice function C_p defined for all O in \mathcal{Q} by

$$C_p(O) := \{u \in O : (\forall v \in O)(E_p(u) \geq E_p(v) \text{ and } u \not\prec v)\}. \quad (11)$$

Proposition 20. *The set of desirable options D_p and the choice function C_p are coherent and compatible, and moreover $C_p = C_{D_p}$.*

This result allows us to introduce the following, second special case of ‘infimum of maximality’ choice functions.

Definition 12. With any non-empty set of mass functions K ,⁵ we associate the corresponding *E-admissible* choice function $C_K^E := \inf\{C_p : p \in K\} = C_{\{D_p : p \in K\}}$.

⁵Although Levi's notion of E-admissibility was originally concerned with *convex closed* sets of mass functions [17, Chapter 5], we impose no such requirement here on the set K .

Proposition 21. *Given any non-empty set of mass functions K , we have for all O in \mathcal{Q} that*

$$C_K^E(O) = \left\{ u \in O : (\exists p \in K) E_p(u) \in \arg \max_{v \in O} E_p(v) \right\} \cap C_v(O).$$

The following proposition establishes a connection between M-admissible and E-admissible choice functions.

Proposition 22. *For any non-empty set of mass functions K , $C_K^E \subseteq C_{\hat{\mathcal{D}}_K}^M$, where $\hat{\mathcal{D}}_K := \bigcup_{p \in K} \hat{\mathcal{D}}_{D_p} \subseteq \hat{\mathcal{D}}$.*

The key for this result is that, for any mass function p , there is in general more than one coherent set of desirable options D that is associated with it by means of the formula

$$E_p(u) = \sup \{ \mu : u - \mu \in D \}. \quad (12)$$

This is why in general coherent sets of desirable options are more informative than (sets of) mass functions as a belief model. Among all the coherent sets of desirable options satisfying Equation (12) with respect to a fixed p , the least informative one is the one given by Equation (10), which is usually referred to as the set of *strictly desirable* gambles associated with p within the imprecise probability literature. This in turn gives rise to a coherent choice function that will be less informative than one determined by a maximal set of options that is compatible with p by means of Equation (12).

Thus, the choice between E-admissible and M-admissible coherent choice functions can be made by considering our attitude towards imprecision, that determines the use of strictly desirable or maximal sets of options: the former are as conservative as possible, and make a choice only when it is implied by the probability mass function p ; while the latter are as informative as it can be considering the axioms of coherence and the probability mass function p .

The following examples show why choice functions are more powerful than sets of desirable options as uncertainty representations, and elucidate the difference between E-admissible and M-admissible choice functions.

Example 5. Consider the situation where you have a coin with two *identical* sides of unknown type:⁶ either both sides are heads (H), or both sides are tails (T). The random variable that represents the outcome of a coin flip assumes a value in the finite possibility space $\mathcal{X} := \{H, T\}$. The options we consider are gambles: real-valued functions on \mathcal{X} , which constitute the two-dimensional vector space $\mathbb{R}^{\mathcal{X}}$, ordered by the pointwise order. We model this situation using (a) coherent sets of desirable options, (b) M-admissible choice functions, and (c) E-admissible choice functions. In all three cases we start from two simple models: one that describes practical certainty of H and another that describes practical certainty of T, and we take their infimum—the most informative model that is still at most as informative as either—as a candidate model for the coin problem.

For (a), we use two coherent sets of desirable options D_H and D_T , expressing practical certainty of H and T, respectively, given by the *maximal* sets of desirable options $D_H := \mathcal{V}_{>0} \cup \{u \in \mathcal{V} : u_H > 0\}$ and $D_T := \mathcal{V}_{>0} \cup \{u \in \mathcal{V} : u_T > 0\}$, where u_H and u_T denote the values of the gamble u in H and T, respectively. The model for the coin with two identical sides is then $D_H \cap D_T = \mathcal{V}_{>0}$. This vacuous model D_v is incapable of distinguishing between this situation and the one where we are completely ignorant about the coin.

⁶The example can be trivially reformulated to consider a random variable taking values in a binary possibility space, with only one of those elements occurring; however we think the use of a coin adds some intuition and helps in particular when we introduce indifference in Section 5 later on.

For an approach (b) that distinguishes between these two situations, we draw inspiration from Proposition 15: instead of working with the sets of desirable options themselves, we move to the corresponding choice functions $C_H := C_{D_H}$ and $C_T := C_{D_T}$, where

$$\begin{aligned} C_H(O) &= \{u \in O : (\forall v \in O)v - u \notin D_H\} = \arg \max\{u_H : u \in O\} \cap C_V(O) \\ &= \arg \max\{v_T : v \in \arg \max\{u_H : u \in O\}\} \\ C_T(O) &= \arg \max\{u_T : u \in O\} \cap C_V(O) \\ &= \arg \max\{v_H : v \in \arg \max\{u_T : u \in O\}\} \end{aligned}$$

for all O in \mathcal{Q} . We infer that $|C_H(O)| = |C_T(O)| = 1$ for every O in \mathcal{Q} : for instance in the case of C_H , note that amongst all the options attaining the maximum value on heads, exactly one of them is undominated. The M-admissible choice function we are looking for is $C_{\{D_H, D_T\}}^M = \inf\{C_H, C_T\}$, which selects at most two options from each option set. It is given by

$$C_{\{D_H, D_T\}}^M(O) = (\arg \max\{u_H : u \in O\} \cup \arg \max\{u_T : u \in O\}) \cap C_V(O)$$

for all O in \mathcal{Q} , and differs from the vacuous choice function C_V . Indeed, consider the particular option set $O = \{u, v, w\}$, where $u = (1, 0)$, $v = (0, 1)$ and $w = (1/2, 1/2)$. Then $C_{\{D_H, D_T\}}^M(O) = \{u, v\} \neq O = C_V(O)$.

For (c), the set of mass functions K consists of the two degenerate mass functions: $K = \{p_H, p_T\}$, where $p_H = (1, 0)$ and $p_T = (0, 1)$. The corresponding expectations $E_H := E_{p_H}$ and $E_T := E_{p_T}$ satisfy $E_H(u) = u_H$ and $E_T(u) = u_T$ for all u in \mathcal{V} . So we see that $C_{p_H} = C_H$ and $C_{p_T} = C_T$, and therefore this approach leads to the same choice function as the previous one: $C_{\{p_H, p_T\}}^E = C_{\{D_H, D_T\}}^M = \inf\{C_H, C_T\}$. \square

The example above shows that the correspondence between desirability and choice functions is not a complete inf-homomorphism, as illustrated also at the end of Section 3.

Example 6. In this example, we illustrate the difference between E-admissible and M-admissible choice functions. We consider the same finite possibility space $\mathcal{X} := \{H, T\}$ as in Example 5, with the same option space and vector ordering. For both E-admissibility and M-admissibility, we each time consider the least informative choice functions: the E-admissible choice function $C_{\Sigma_{\mathcal{X}}}^E$ associated with set of all mass functions $K = \Sigma_{\mathcal{X}}$, and the M-admissible choice function $C_{D_V}^M$ associated with the set of all maximal sets of desirable options $\hat{D}_{D_V} = \hat{D}$. Since $C_{\Sigma_{\mathcal{X}}}^E$ and $C_{D_V}^M$ are the most conservative E-admissible, respectively M-admissible choice functions, we wonder about their relation, as well as their relation with the vacuous choice function C_V . We find that $C_{\Sigma_{\mathcal{X}}}^E \subseteq C_{D_V}^M$. Indeed, consider any O in \mathcal{Q} and any u in $C_{D_V}^M(O)$, being equivalent to $0 \in C_{D_V}^M(O - \{u\})$ by Proposition 3(i), whence

$$u \in C_{D_V}^M(O) \Leftrightarrow (\exists \hat{D} \in \hat{D})(\forall v \in O - \{u\})v \notin \hat{D} \Leftrightarrow (\exists \hat{D} \in \hat{D})O - \{u\} \cap \hat{D} = \emptyset.$$

Since for every \hat{D} in \hat{D} , there is some mass function p in $\Sigma_{\mathcal{X}}$ such that $D_p \subseteq \hat{D}$ [it suffices to consider the mass function p corresponding with \hat{D}], we find that

$$u \in C_{D_V}^M(O) \Rightarrow (\exists p \in \Sigma_{\mathcal{X}})O - \{u\} \cap D_p = \emptyset \Leftrightarrow (\exists p \in \Sigma_{\mathcal{X}})u \in C_{D_p}(O).$$

or in other words, $u \in C_{D_V}^M(O)$ implies that $u \in C_{\Sigma_{\mathcal{X}}}^E(O)$, whence $C_{D_V}^M(O) \subseteq C_{\Sigma_{\mathcal{X}}}^E(O)$. By the definition of the vacuous choice function, we have as an intermediate result that $C_V \subseteq C_{\Sigma_{\mathcal{X}}}^E \subseteq C_{D_V}^M$.

Both inequalities are strict; to show that $C_V \neq C_{\Sigma_{\mathcal{X}}}^E$, consider the option set $O := \{0, u, v\}$, where $u = (1, -1/4)$ and $v = (-1/4, 1)$. Because all options in O are pointwise undominated in

O , we find that $C_v(O) = O$, and in particular, that $0 \in C_v(O)$. On the other hand, it follows from Proposition 21 that

$$0 \in C_{\Sigma_{\mathcal{X}}}^E(O) \Leftrightarrow (\exists p \in \Sigma_{\mathcal{X}})(0 \geq p(H) - \frac{1}{4}p(T) \text{ and } 0 \geq -\frac{1}{4}p(H) + p(T)),$$

which would imply that $(\exists p \in \Sigma_{\mathcal{X}})(0 \geq \frac{3}{4}p(H) + \frac{3}{4}p(T))$, which is impossible.

More importantly, we also have that $C_{\Sigma_{\mathcal{X}}}^E \neq C_{D_v}^M$. Consider the option set $O := \{0, u, -u\}$, where $u = (1, -1)$. Then for the specific mass function $p := (1/2, 1/2) \in \Sigma_{\mathcal{X}}$, we find that $0 \in C_p(O)$, whence $0 \in C_{\Sigma_{\mathcal{X}}}^E(O)$. To show that $0 \in C_p^E(O)$, infer that $0 = E_p(0) = E_p(u) = E_p(-u)$, and use Proposition 21 as a characterisation for the E-admissible choice functions. On the other hand, $0 \in C_{D_v}^M(O)$ is equivalent to $u, -u \notin \hat{D}$ for some \hat{D} in $\hat{\mathcal{D}}$. But $u \notin \hat{D}$ and $-u \notin \hat{D}$ implies that $-u \in \hat{D}$ and $u \in \hat{D}$ by Proposition 10, a contradiction. So $0 \notin C_{D_v}^M(O)$, whence $C_{\Sigma_{\mathcal{X}}}^E \neq C_{D_v}^M$.

This example shows that $C_v = C_{\mathcal{D}} \subset C_{\Sigma_{\mathcal{X}}}^E \subset C_{\mathcal{D}} = C_{D_v}^M$. \square

As remarked by a reviewer, we can interpret the example above in terms of the choice relation associated with a choice function, in the following manner: in the case of a complete preference relation we always have that $\{0\} \prec_C \{u, -u\}$ for every option u . This is not the case for those induced by sets of strictly desirable options, such as the coherent choice function C_p in the example, which therefore cannot be obtained as infima of a family of complete choice relations (as are those given by M -admissibility).

In Section 5.7, we will present another example that illustrates the difference between E-admissible and M-admissible choice functions.

There are other coherent sets of desirable options that can be associated with a mass function p and that are intermediate between the strictly desirable and the maximal ones. One example are the so-called *lexicographic* sets of desirable options, which we have investigated in detail elsewhere [32]. Taking into account Proposition 18, we can also define coherent choice functions by taking the infimum of a family of coherent choice functions determined by such lexicographic sets. Indeed, this provides another example of coherent choice function that admits an axiomatic characterisation in some cases, as reported in [32].

To conclude this section, we want to mention that there are other popular choice rules besides maximality and E-admissibility, such as, amongst others, Γ -*maximin*, Γ -*maximax* and *interval dominance* [30]. However, they are not coherent: none of them is guaranteed to satisfy, amongst others, Axiom C_{4b}.

5. INDIFFERENCE

5.1. Indifference and desirability. For sets of desirable gambles, there is a systematic way of modelling indifference [8, 11, 20]. In what follows, we recall how to do this, but formulate everything in terms of the more abstract notion of options, rather than gambles.

In addition to a subject's set of desirable options D —the options he strictly prefers to the zero option—we can also consider the options that he considers to be *equivalent* to the zero option. We call these options *indifferent*. A set of indifferent options I is simply a subset of \mathcal{V} , but as before with desirable options, we pay special attention to *coherent* sets of indifferent options.

Definition 13. A set of indifferent options I is called coherent if for all u, v in \mathcal{V} and λ in \mathbb{R} :

- I₁. $0 \in I$;
- I₂. if $u \in \mathcal{V}_{>0} \cup \mathcal{V}_{<0}$ then $u \notin I$;
- I₃. if $u \in I$ then $\lambda u \in I$;

I₄. if $u, v \in I$ then $u + v \in I$.

Taken together, Axioms I₃ and I₄ are equivalent to imposing that $\text{span}(I) = I$, and due to Axiom I₁, I is non-empty and therefore a linear subspace of \mathcal{V} .

The interaction between indifferent and desirable options is subject to rationality criteria as well: they should be compatible with one another.

Definition 14. Given a set of desirable options D and a coherent set of indifferent options I , we call D *compatible* with I if $D + I \subseteq D$.

The smallest such compatible coherent set of indifferent options is always the trivial one, given by the null space $I = \{0\}$. The idea behind Definition 14 is that adding an indifferent option to a desirable option does not make it non-desirable, or also, that adding a desirable option to an indifferent one makes it desirable.

Since $D \subseteq D + I$ due to Axiom I₁, compatibility of D and I is equivalent to $D + I = D$. An immediate consequence of compatibility between a coherent set of desirable options D and a coherent set of indifferent options I is that $D \cap I = \emptyset$, meaning that no option can be assessed as desirable—strictly preferred to the zero option—and indifferent—equivalent to the zero option—at the same time. To see this, if *ex absurdo* $D \cap I \neq \emptyset$, then there is some u in D such that $u \in I$. But then, by Axiom I₃, also $-u \in I$, so compatibility of D with I implies that $u + (-u) = 0 \in D$, a contradiction with Axiom D₁.

5.2. Indifference and quotient spaces. In order to introduce indifference for choice functions, we build on a coherent set of indifferent options I , as defined in Definition 13. Two options u and v are considered to be indifferent, to a subject, whenever $v - u$ is indifferent to the zero option, or in other words whenever $v - u \in I$. The underlying idea will be that we identify indifferent options, and choose between equivalence classes of indifferent options, rather than between single options. We formalise this intuition below.

We collect all options that are indifferent to an option $u \in \mathcal{V}$ into the *equivalence class*

$$[u] := \{v \in \mathcal{V} : v - u \in I\} = \{u\} + I.$$

We also denote $[u]$ as u/I . Of course, $[0] = \{0\} + I = I$ is a linear subspace, and the $[u] = \{u\} + I$ are affine subspaces of \mathcal{V} . The set of all these equivalence classes is the *quotient space*

$$\mathcal{V}/I := \{[u] : u \in \mathcal{V}\} = \{\{u\} + I : u \in \mathcal{V}\} = \{u/I : u \in \mathcal{V}\}.$$

This quotient space is a vector space under the vector addition, given by

$$[u] + [v] = \{u\} + I + \{v\} + I = \{u + v\} + I = [u + v] \text{ for } u, v \in \mathcal{V},$$

and the scalar multiplication, given by

$$\lambda[u] = \lambda(\{u\} + I) = \{\lambda u\} + I = [\lambda u] \text{ for } u \in \mathcal{V} \text{ and } \lambda \in \mathbb{R}.$$

$[0] = I$ is the additive identity of \mathcal{V}/I .

That we identify indifferent options, and therefore express preferences between equivalence classes of indifferent options, essentially means that we define choice functions on $\mathcal{Q}(\mathcal{V}/I)$. But in order to characterise coherence for such choice functions, we need to introduce a convenient vector ordering on \mathcal{V}/I that is appropriately related to the vector ordering on \mathcal{V} ; see Section 2.2. For two elements $[u]$ and $[v]$ of \mathcal{V}/I , we define

$$[u] \leq [v] \Leftrightarrow (\exists w \in I) u \leq v + w, \tag{13}$$

and as usual, the strict variant of the vector ordering on \mathcal{V}/I is characterised by

$$[u] < [v] \Leftrightarrow ([u] \leq [v] \text{ and } [u] \neq [v]).$$

We begin by showing that this is indeed a vector ordering:

Proposition 23. *The ordering \leq on \mathcal{V}/I is a vector ordering, and $[u] < [v] \Leftrightarrow (\exists w \in I)u < v + w$ for any u, v in \mathcal{V} .*

We use the notation $O/I := \{[u] : u \in O\} = \{u/I : u \in O\}$ for the option set of equivalence classes $[u]$ associated with the options u in an option set O in $\mathcal{Q}(\mathcal{V})$. The map \cdot/I is an onto map from $\mathcal{Q}(\mathcal{V})$ to $\mathcal{Q}(\mathcal{V}/I)$ that preserves set inclusion:

$$(\forall O_1, O_2 \in \mathcal{Q}(\mathcal{V})) (O_1 \subseteq O_2 \Rightarrow O_1/I \subseteq O_2/I). \quad (14)$$

5.3. Quotient spaces and sets of desirable options. We use this quotient space to prove interesting characterisations of indifference for sets of desirable options.

Proposition 24. *A set of desirable options $D \subseteq \mathcal{V}$ is compatible with a coherent set of indifferent options I if and only if there is some (representing) set of desirable options $D' \subseteq \mathcal{V}/I$ such that $D = \{u : [u] \in D'\} = \bigcup D'$. Moreover, the representing set of desirable options is unique and given by $D' = D/I := \{[u] : u \in D\}$.*

This, together with the definition of compatibility, shows that the correspondence between sets of desirable options on \mathcal{V} and (their representing) sets of desirable options on \mathcal{V}/I is one-to-one and onto. It also preserves coherence.

Proposition 25. *Consider any set of desirable options $D \subseteq \mathcal{V}$ that is compatible with a coherent set of indifferent options I , and its representing set of desirable options $D/I \subseteq \mathcal{V}/I$. Then D is coherent if and only if D/I is.*

5.4. Quotient spaces and choice functions. The discussion above inspires us to combine indifference with choice functions in the following manner: given a coherent set of indifferent options I , we say that a choice function is compatible with it when it is determined by its restriction to the quotient space that I induces.

Definition 15. We call a choice function C on $\mathcal{Q}(\mathcal{V})$ *compatible* with a coherent set of indifferent options I if there is some *representing* choice function C' on $\mathcal{Q}(\mathcal{V}/I)$ such that $C(O) = \{u \in O : [u] \in C'(O/I)\}$ for all O in $\mathcal{Q}(\mathcal{V})$.

This definition can be stated in terms of choice relations as well. We call a choice relation \triangleleft on $\mathcal{Q}(\mathcal{V})$ compatible with I if there is some representing choice relation \triangleleft' on $\mathcal{Q}(\mathcal{V}/I)$ such that $O_1 \triangleleft O_2 \Leftrightarrow O_1/I \triangleleft' O_2/I$ for all O_1 and O_2 in $\mathcal{Q}(\mathcal{V})$.

Proposition 26. *Consider any choice function C that is compatible with I , then \triangleleft_C is compatible with I as well. Conversely, consider any choice relation \triangleleft that is compatible with I , then C_\triangleleft is compatible with I as well.*

Definition 15 allows for characterisations that are similar to the ones for desirability in Propositions 24 and 25. If a choice function on $\mathcal{Q}(\mathcal{V})$ is compatible with I then the representing choice function on $\mathcal{Q}(\mathcal{V}/I)$ is necessarily unique, and we denote it by C/I :

Proposition 27. *For any choice function C on $\mathcal{Q}(\mathcal{V})$ that is compatible with some coherent set of indifferent options I , the unique representing choice function C/I on $\mathcal{Q}(\mathcal{V}/I)$ is given by $C/I(O/I) := C(O)/I$ for all O in $\mathcal{Q}(\mathcal{V})$. Hence also*

$$C(O) = O \cap \left(\bigcup C/I(O/I) \right) \text{ for all } O \text{ in } \mathcal{Q}(\mathcal{V}).$$

This, together with the definition of compatibility, shows that the correspondence between choice functions on $\mathcal{Q}(\mathcal{V})$ and (their representing) choice functions on $\mathcal{Q}(\mathcal{V}/I)$ is one-to-one and onto. It also preserves coherence:

Proposition 28. *Consider any choice function C on $\mathcal{Q}(\mathcal{V})$ that is compatible with a coherent set of indifferent options I , and its representing choice function C/I on $\mathcal{Q}(\mathcal{V}/I)$. Then C is coherent if and only if C/I is.*

To conclude this general discussion of indifference for choice functions, we mention that it is closed under arbitrary infima, which enables *conservative inference under indifference*: we can consider the least informative choice function that is compatible with some assessments and is still compatible with a given coherent set of indifferent options.

Proposition 29. *Consider any coherent set of indifferent options I , and any non-empty collection of coherent choice functions $\{C_i : i \in \mathcal{I}\}$ that are compatible with I , then its coherent infimum $\inf\{C_i : i \in \mathcal{I}\}$ is compatible with I as well, and $C/I = \inf\{C_i/I : i \in \mathcal{I}\}$.*

In particular, the least informative coherent choice function C that is compatible with a coherent set of indifferent options I corresponds to the case where the associated choice function C/I on the quotient space $\mathcal{Q}(\mathcal{V}/I)$ is the vacuous one C_v , meaning that

$$C(O) = \{u \in O : [u] \in C_v(O/I)\} = \{u \in O : (\forall v \in O)[u] \not\prec [v]\} \text{ for all } O \text{ in } \mathcal{Q}(\mathcal{V}).$$

5.5. The relation with other definitions of indifference. Seidenfeld [24] has a rather different approach to combining a notion of indifference with choice functions. After making the necessary translation from horse lotteries to gambles, it can be summarised as follows. Rather than, as we have done above, starting out with a notion of indifference and then looking at which choice functions are compatible with it, he starts from a given choice function C , and associates a binary relation \approx_C on \mathcal{V} with it as follows:

$$u \approx_C v \Leftrightarrow (\forall O \in \mathcal{Q})(\{u, v\} \subseteq O \Rightarrow (u \in C(O) \Leftrightarrow v \in C(O))) \text{ for all } u, v \text{ in } \mathcal{V}. \quad (15)$$

The idea behind this definition is that two options are considered to be related if both options are either chosen or rejected, whenever both are available. This relation has the following interesting properties, which are instrumental in elucidating the relationship between Seidenfeld's approach and ours:

Proposition 30. *Let C be a coherent choice function. Then \approx_C is an equivalence relation that furthermore satisfies*

- (i) $u \approx_C v \Rightarrow \lambda u \approx_C \lambda v$;
- (ii) $u \approx_C v \Rightarrow u + w \approx_C v + w$;

for all $u, v, w \in \mathcal{V}$ and all $\lambda \in \mathbb{R}$. As a consequence, the set $I_C := \{u \in \mathcal{V} : u \approx_C 0\}$ is a coherent set of indifferent options.

Indeed, the coherent set of indifferent options I_C turns out to be the largest that C is compatible with:

Proposition 31. *Consider a coherent set of indifferent options I and a coherent choice function C on $\mathcal{Q}(\mathcal{V})$. Then C is compatible with I if and only if I is a linear subspace of I_C .*

It follows that if C is compatible with I , then C is automatically also compatible with any subspace of I , and that the largest linear subspace that C is compatible with, is I_C . Also, this proposition shows that the smallest set of indifferent option that is compatible with C , is given by $I = \{0\}$, and that set of indifferent options is compatible with any coherent choice function. Seidenfeld's approach starts from a choice function, and identifies the coarsest equivalence—or indifference—relation that is compatible with it. Though we have seen that it is related, our approach, because it starts out with an indifference relation, goes the other way around, is more constructive, and is better suited for studying which choice functions are compatible with a given indifference relation, and therefore also for

performing conservative inference for choice functions under indifference. We will come back to this idea in Sections 5.7 and 5.8.

Proposition 31 also shows that our approach is closely connected to Bradley’s [4, Section 2], which he has linked to the notion of sequential coherence: for him, for a choice function to ‘*satisfy*’—in our words: to ‘be compatible with’—an equivalence relation, it must be finer than Seidenfeld’s \approx_C .

Finally, note that a relation analogous to Equation (15) can be established for coherent sets of desirable options D : we simply define $u \approx_D v \Leftrightarrow u \approx_{C_D} v$ for all u and v in \mathcal{V} . Then

$$\begin{aligned} u \approx_D v &\Leftrightarrow (\forall O \in \mathcal{Q}) (\{u, v\} \subseteq O \Rightarrow (0 \in C_D(O - \{u\}) \Leftrightarrow 0 \in C_D(O - \{v\}))) \\ &\Leftrightarrow (\forall O \in \mathcal{Q}) (\{u, v\} \subseteq O \Rightarrow (O - \{u\} \cap D \neq \emptyset \Leftrightarrow O - \{v\} \cap D \neq \emptyset)) \text{ for all } u, v \text{ in } \mathcal{V}, \end{aligned}$$

using Axiom C4b for the first equivalence and Equation (8) for the second one. Use Lemma 32 to infer that

$$u \approx_D v \Leftrightarrow D + \{u\} = D + \{v\} \text{ for all } u, v \text{ in } \mathcal{V}.$$

Lemma 32. *Consider any coherent set of desirable options D . Then*

$$(\forall O \in \mathcal{Q}) (\{u, v\} \subseteq O \Rightarrow (O - \{u\} \cap D \neq \emptyset \Leftrightarrow O - \{v\} \cap D \neq \emptyset)) \Leftrightarrow D + \{u\} = D + \{v\}$$

for all u and v in \mathcal{V} .

5.6. The relation with desirability. There is an interesting relationship between the coherent choice functions and the coherent sets of desirable options that are compatible with a fixed coherent set of indifferent options.

Proposition 33. *Consider any coherent set of indifferent options I .*

- (i) *If C is any coherent choice function compatible with I , then the corresponding coherent set of desirable options D_C is also compatible with I , and $D_C/I = D_C|I$.*
- (ii) *If D is any coherent set of desirable options compatible with I , then the corresponding coherent choice function C_D is also compatible with I , and $C_D/I = C_D|I$.*

5.7. Example: fair coins. To exhibit the power and simplicity of our definition of indifference, we reconsider the finite possibility space $\mathcal{X} := \{H, T\}$ of Example 5, where the vector space \mathcal{V} is again the two-dimensional vector space $\mathbb{R}^{\mathcal{X}}$ of real-valued functions on \mathcal{X} , or gambles, and the vector ordering \leq is the usual pointwise ordering of gambles.

We want to express indifference between heads and tails, or in other words between \mathbb{I}_H and \mathbb{I}_T , where $\mathbb{I}_H := (1, 0)$ and $\mathbb{I}_T := (0, 1)$. This means that $\mathbb{I}_H - \mathbb{I}_T$ is considered equivalent to the zero gamble, so the linear space of all gambles that are equivalent to zero—or in other words, the set of indifferent gambles (or options)—is then given by

$$I = \{\lambda(\mathbb{I}_H - \mathbb{I}_T) : \lambda \in \mathbb{R}\} = \{u \in \mathbb{R}^{\mathcal{X}} : E_p(u) = 0\},$$

where E_p is the expectation associated with the uniform mass function $p = (1/2, 1/2)$ on $\{H, T\}$, associated with a fair coin: $E_p(u) := \frac{1}{2}[u_H + u_T]$. So, for any option u in $\mathbb{R}^{\mathcal{X}}$ —any real-valued function on \mathcal{X} :

$$[u] = \{u\} + I = \{v \in \mathbb{R}^{\mathcal{X}} : E_p(v) = E_p(u)\},$$

which tells us that the equivalence class $[u]$ can be characterised by the common uniform expectation $E_p(u)$ of its elements. Therefore, $\mathbb{R}^{\mathcal{X}}/I$ has unit dimension, and we can identify it with the real line \mathbb{R} . The vector ordering between equivalence classes is given by, using Equation (13):

$$[u] \leq [v] \Leftrightarrow (\exists \lambda \in \mathbb{R}) u \leq v + \lambda(\mathbb{I}_H - \mathbb{I}_T)$$

$$\begin{aligned}
&\Leftrightarrow (\exists \lambda \in \mathbb{R})(u_H \leq v_H + \lambda \text{ and } u_T \leq v_T - \lambda) \\
&\Leftrightarrow (\exists \lambda \in \mathbb{R})u_H - v_H \leq \lambda \leq -u_T + v_T \\
&\Leftrightarrow u_H - v_H \leq -u_T + v_T \Leftrightarrow E_p(u) \leq E_p(v),
\end{aligned}$$

and similarly $[u] < [v] \Leftrightarrow E_p(u) < E_p(v)$ for all u, v in $\mathbb{R}^{\mathcal{X}}$. Hence, the strict vector ordering $<$ on $\mathbb{R}^{\mathcal{X}}/I$ is total, so we infer from the argumentation in Example 3 that there is only one representing choice function, namely the vacuous one. Therefore, there is only one choice function C on $\mathcal{Q}(\mathbb{R}^{\mathcal{X}})$ that is compatible with I , namely, the one that has the vacuous choice function C_v on $\mathcal{Q}(\mathbb{R}^{\mathcal{X}}/I)$ as its representation C/I . Recall that for any O in $\mathcal{Q}(\mathbb{R}^{\mathcal{X}})$:

$$\begin{aligned}
C_v(O/I) &= \{[u] : (\forall [v] \in O/I)[u] \neq [v]\} \\
&= \{[u] : (\forall [v] \in O/I)[v] \leq [u]\} \\
&= \{[u] : (\forall [v] \in O/I)E_p(v) \leq E_p(u)\},
\end{aligned}$$

and therefore

$$C(O) := \{u \in O : (\forall v \in O)E_p(v) \leq E_p(u)\} = C_{\{p\}}^E(O).$$

The indifference assessment between heads and tails leaves us no choice but to use an E-admissible model for a probability mass function, associated with a fair coin.

The choice function C is therefore based on E-admissibility, but is not compatible with M-admissibility. To see this, consider the set of options $O := \{w, 0, -w\}$ with $w := (1, -1)$, so $w_H + w_T = 0$. Hence $C(O) = O$; but no M-admissible choice function will select 0 in O : observe that $0 \notin C_{\hat{D}}(O)$ for all $\hat{D} \in \mathcal{D}'$, because $0 \in C_{\hat{D}}(O)$ would imply that $\{w, -w\} \cap \hat{D} = \emptyset$, contradicting that \hat{D} is a maximal set of desirable options by Proposition 10.

5.8. Choice functions and symmetry. As another example showing how powerful our approach to dealing with choice and indifference is, we will prove a simple and elegant representation result that tells us how to perform conservative inference with choice functions under a permutation symmetry assessment.

We consider a finite possibility space \mathcal{X} , where the vector space \mathcal{V} of options is the finite-dimensional vector space $\mathbb{R}^{r, \mathcal{X}}$, of \mathbb{R}^r -valued functions on \mathcal{X} , or *vector-valued gambles* on the outcome of an uncertain variable X in \mathcal{X} . The vector ordering \leq is the usual pointwise ordering of such vector-valued gambles.⁷

We assume there is symmetry lurking behind the uncertain variable X , represented by a group \mathcal{P} of permutations of the set of possible outcomes \mathcal{X} —the idea being that a subject assesses that no distinction should be made between an outcome x and its permutations πx , for $\pi \in \mathcal{P}$ —or in other words, between the variable X and its permutations πX . If we consider any vector-valued gamble $u(X)$ on the variable X , then the subject will therefore be indifferent between the uncertain vector-valued rewards $u(X)$ and $u(\pi X)$. The smallest coherent⁸ set of indifferent options $I_{\mathcal{P}}$ that corresponds to this indifference assessment, is therefore given by

$$I_{\mathcal{P}} := \text{span}(\{u - \pi^t u : u \in \mathcal{V} \text{ and } \pi \in \mathcal{P}\}),$$

⁷The reason why we work with vector-valued gambles here rather than real-valued ones, is because Zaffalon and Miranda [36] have shown that this leads to an approach to modelling uncertainty that is even more general than the typical imprecise probability approach.

⁸The requirement that $I_{\mathcal{P}} \cap \mathcal{V}_{<0} = \emptyset$ —or equivalently $I_{\mathcal{P}} \cap \mathcal{V}_{>0} = \emptyset$ —is related to the left amenability of the finite permutation group \mathcal{P} [9, 34], and is easily shown to be satisfied.

where we defined the linear permutation operator π^t on the linear space of options (vector-valued gambles) \mathcal{V} by $\pi^t u = u \circ \pi$, or in other words

$$(\pi^t u)(x) := u(\pi x) \text{ for all } u \text{ in } \mathcal{V}, x \text{ in } \mathcal{X} \text{ and } \pi \text{ in } \mathcal{P}.$$

Let us, for any x in \mathcal{X} , define the *permutation invariant atom* $[x]_{\mathcal{P}}$ containing x as

$$[x]_{\mathcal{P}} := \{\pi x : \pi \in \mathcal{P}\}.$$

These atoms constitute a partition of \mathcal{X} , and we denote the set of all of them by $\mathcal{A}_{\mathcal{P}} := \{[x]_{\mathcal{P}} : x \in \mathcal{X}\}$. A vector-valued gamble u is called \mathcal{P} -invariant if $\pi^t u = u$ for all π in \mathcal{P} , and it is not hard to see that this is equivalent to u being constant on the invariant atoms. The set of all \mathcal{P} -invariant vector-valued gambles is denoted by $\mathcal{V}_{\mathcal{P}}$, and it is a linear subspace of \mathcal{V} that is clearly isomorphic to the linear space of all vector-valued functions on $\mathcal{A}_{\mathcal{P}}$, whose dimension $r|\mathcal{A}_{\mathcal{P}}|$ is typically much lower than that of \mathcal{V} .

A choice function that takes the symmetry assessment into account is—as we have argued—one that is compatible with $I_{\mathcal{P}}$ and all of its subspaces. What we will do now, is to investigate how such compatible choice functions can be represented by choice functions on a typically much lower-dimensional option space: symmetry reduces complexity. Most of the work for this has already been done in Definition 15 and Proposition 27, which indeed states that choice functions compatible with $I_{\mathcal{P}}$ can be represented uniquely by choice functions on the lower-dimensional quotient space $\mathcal{V}/I_{\mathcal{P}}$. The only thing we need to do, is take a closer look at this quotient space and its elements.

Let us, therefore, define the transformation $\text{inv}_{\mathcal{P}}$ on \mathcal{V} as follows:

$$\text{inv}_{\mathcal{P}} u := \frac{1}{|\mathcal{P}|} \sum_{\pi \in \mathcal{P}} \pi^t u, \text{ for all } u \text{ in } \mathcal{V}. \quad (16)$$

It satisfies the following very interesting properties:

Proposition 34. *inv_℘ is a linear transformation of ℳ, and*

- (i) $\text{inv}_{\mathcal{P}} \circ \pi^t = \text{inv}_{\mathcal{P}} = \pi^t \circ \text{inv}_{\mathcal{P}}$ for all π in \mathcal{P} ;
- (ii) $\text{inv}_{\mathcal{P}} \circ \text{inv}_{\mathcal{P}} = \text{inv}_{\mathcal{P}}$;
- (iii) $\text{kern}(\text{inv}_{\mathcal{P}}) = I_{\mathcal{P}}$;
- (iv) $\text{rng}(\text{inv}_{\mathcal{P}}) = \mathcal{V}_{\mathcal{P}}$.

Moreover, for any u, v in \mathcal{V} , we have that $v \in u/I_{\mathcal{P}} \Leftrightarrow \text{inv}_{\mathcal{P}} v = \text{inv}_{\mathcal{P}} u$,

The various statements in this proposition tell us that $\text{inv}_{\mathcal{P}}$ is a linear projection operator that maps any vector-valued gamble u to the corresponding uniquely \mathcal{P} -invariant member $\text{inv}_{\mathcal{P}} u$ of the equivalence class $u/I_{\mathcal{P}}$, which is essentially a vector-valued gamble on $\mathcal{A}_{\mathcal{P}}$.

By Proposition 27, every coherent choice function C on \mathcal{V} that is compatible with $I_{\mathcal{P}}$ therefore has a unique representing coherent choice function $C_{\mathcal{P}}$ on the typically much lower-dimensional linear space of all vector-valued gambles on $\mathcal{A}_{\mathcal{P}}$, with

$$C(O) = \{u \in O : \text{inv}_{\mathcal{P}} u \in C_{\mathcal{P}}(\text{inv}_{\mathcal{P}} O)\}.$$

6. CONCLUSIONS

In this paper, we have shown how the theory of choice functions may be used, with a choice of axioms to model decision problems where the choices are not uniquely determined by the pairwise comparisons, and where models of sets of desirable gambles therefore come up short.

We have made our formalisation in terms of sets of options, because this more general structure has allowed us to make a connection with modelling indifference. Nevertheless, our models remain applicable in particular to choices on sets of gambles, and thus encompass

much of the behavioural theory of imprecise probabilities as a particular case. A deeper study of the definition in terms of horse lotteries, which is the model considered by Seidenfeld [24], is one of our future lines of research. We have established that choice functions can be equivalently formulated as choice relations, that is, binary relations in terms of sets of options. This formulation adds intuition to some of the results in the paper.

We have seen that three important types of decision models can be subsumed into our theory of choice functions: one, as we have said, is that of sets of desirable options, which give rise to the choice functions defined via maximality; the others are those based on a family of sets of probability measures or, closely related, a family of maximal sets of desirable gambles. We have named these two types of models E-admissible and M-admissible choice functions, and we have shown that there are subtle differences between them. Related to this, one future line of research would be the study of *maximal* coherent choice functions, which are those that are not above any other choice function. As suggested by one reviewer, it may be useful here to use the equivalent formulation in terms of choice relations and to use the notion of *consequence* operator [6].

One of the advantages of sets of desirable options over sets of probability measures is that they are more suited to conditioning when sets of (lower) probability zero are involved, as discussed for instance by Quaeghebeur [19]. We intend to investigate in detail in future work how conditioning for choice functions should be defined.

Finally, we have given an axiomatisation of sets of indifferent options and their compatibility with choice functions, and have shown how our definition is related to other approaches proposed in the literature. The investigation of this interesting connection under our framework is another open avenue for future research.

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APPENDIX A. PROOFS

Proof of Proposition 1. We have to check that \triangleleft_C satisfies the two conditions of Definition 2. For (i), consider any O , O_1 and O_2 in \mathcal{Q} such that $O_2 \setminus O_1 \subseteq O \subseteq O_1 \cup O_2$. Then $O_1 \cup O = O_1 \cup O_2$, which sets off the following cascade of equivalences:

$$\begin{aligned} O_1 \triangleleft_C O &\Leftrightarrow C(O_1 \cup O) \subseteq O \setminus O_1 \\ &\Leftrightarrow C(O_1 \cup O_2) \subseteq O \setminus O_1 \\ &\Leftrightarrow C(O_1 \cup O_2) \subseteq (O_1 \cup O_2) \cap (O \setminus O_1) \quad \text{because } C(O_1 \cup O_2) \subseteq O_1 \cup O_2 \\ &\Leftrightarrow C(O_1 \cup O_2) \subseteq O_2 \setminus O_1 \Leftrightarrow O_1 \triangleleft_C O_2, \end{aligned}$$

where the penultimate equivalence follows from the equalities: $(O_1 \cup O_2) \cap (O \setminus O_1) = O_2 \cap O \cap O_1^c = O \cap (O_2 \setminus O_1) = O_2 \setminus O_1$, because by assumption $O \supseteq O_2 \setminus O_1$. For (ii), consider any O , O_1 and O_2 in \mathcal{Q} such that $O_1 \cup O_2 \subseteq O$. Then $O_1 \cup O = O_2 \cup O = O$, which sets off the following cascade of equivalences:

$$O_1 \triangleleft_C O \text{ and } O_2 \triangleleft_C O \Leftrightarrow C(O_1 \cup O) \subseteq O \setminus O_1 \text{ and } C(O_2 \cup O) \subseteq O \setminus O_2$$

$$\begin{aligned} &\Leftrightarrow C(O) \subseteq O \setminus O_1 \text{ and } C(O) \subseteq O \setminus O_2 \\ &\Leftrightarrow C(O) \subseteq O \setminus (O_1 \cup O_2) \Leftrightarrow O_1 \cup O_2 \triangleleft_C O. \quad \square \end{aligned}$$

Proof of Proposition 2. For the first statement, consider any choice function C , and any O in \mathcal{Q} , and infer that

$$\begin{aligned} C_{\triangleleft_C}(O) &= \bigcap \{O' \subseteq O : O \setminus O' \triangleleft_C O\} = \bigcap \{O' \subseteq O : C((O \setminus O') \cup O) \subseteq O \setminus (O \setminus O')\} \\ &= \bigcap \{O' \subseteq O : C(O) \subseteq O'\} = C(O), \end{aligned}$$

using the fact that $C(O) \subseteq O$.

For the second statement, consider any choice relation \triangleleft , and any O_1 and O_2 in \mathcal{Q} , and infer that

$$\begin{aligned} O_1 \triangleleft_{C_{\triangleleft}} O_2 &\Leftrightarrow O_1 \cap C_{\triangleleft}(O_1 \cup O_2) = \emptyset \\ &\Leftrightarrow O_1 \cap (O_1 \cup O_2 \setminus \bigcup \{O' \subseteq O_1 \cup O_2 : O' \triangleleft O_1 \cup O_2\}) = \emptyset \\ &\Leftrightarrow O_1 \subseteq \bigcup \{O' \subseteq O_1 \cup O_2 : O' \triangleleft O_1 \cup O_2\}. \end{aligned}$$

Call $O_3 := \bigcup \{O' \subseteq O_1 \cup O_2 : O' \triangleleft O_1 \cup O_2\}$. Since $O_1 \cup O_2$ is finite, and therefore O_3 is a finite union, we infer from Definition 2(ii) that $O_3 \triangleleft O_1 \cup O_2$. If $O_1 \subseteq O_3$ then Definition 2(ii) implies that then $O_1 \triangleleft O_1 \cup O_2$. Conversely, if $O_1 \triangleleft O_1 \cup O_2$, then $O_1 \subseteq O_3$ by the definition of O_3 , so $O_1 \triangleleft_{C_{\triangleleft}} O_2$ is equivalent to $O_1 \triangleleft O_1 \cup O_2$, which is in turn equivalent to $O_1 \triangleleft O_2$ by Definition 2(i).

For the third statement, assume first that $\triangleleft_C = \triangleleft$. Then the corresponding choice functions are equal: $C_{\triangleleft_C} = C_{\triangleleft}$. But we have just shown that $C_{\triangleleft_C} = C$, whence indeed $C = C_{\triangleleft}$. Conversely, assume that $C_{\triangleleft} = C$. Then the corresponding choice relations are equal: $\triangleleft_{C_{\triangleleft}} = \triangleleft_C$. But we have just shown that $\triangleleft_{C_{\triangleleft}} = \triangleleft$, whence indeed $\triangleleft = \triangleleft_C$. \square

Lemma 35. $\tilde{O} \cap C(O) = \emptyset$ is equivalent to $C(O) \subseteq O \setminus \tilde{O}$ for any C in \mathcal{C} and O, \tilde{O} in \mathcal{Q} .

Proof. Since for any choice function it holds that $C(O) \subseteq O$, the inclusion $C(O) \subseteq O \setminus \tilde{O} = O \cap \tilde{O}^c$ is equivalent to $C(O) \subseteq \tilde{O}^c$, that is, to $C(O) \cap \tilde{O} = \emptyset$. \square

Proof of Axiom C_{3a} \Leftrightarrow Equation (3), and of Axiom C_{3b} \Leftrightarrow Equation (4). That Axiom C_{3a} and Equation (3) are equivalent follows immediately from Lemma 35.

Let us establish the equivalence between Axiom C_{3b} and Equation (4). For necessity, consider any O, O_1, O_2 in \mathcal{Q} such that $O \subseteq O_1$ and $O_1 \cap C(O_2) = \emptyset$, or equivalently, $C(O_2) \subseteq O_2 \setminus O_1$, by Lemma 35. Let $\tilde{O} := O \cap O_2$, $\tilde{O}_1 := O_2 \setminus O_1$ and $\tilde{O}_2 := O_2$, then $C(\tilde{O}_2) \subseteq \tilde{O}_1$, and $\tilde{O}_2 \setminus \tilde{O}_1 = O_2 \setminus (O_2 \setminus O_1) = O_1 \cap O_2$, whence $\tilde{O} = O \cap O_2 \subseteq O_1 \cap O_2 = \tilde{O}_2 \setminus \tilde{O}_1$. Due to Axiom C_{3b}, then $C(\tilde{O}_2 \setminus \tilde{O}) \subseteq \tilde{O}_1$, whence $C(O_2 \setminus O) \subseteq O_2 \setminus O_1$. By Definition 1, we have that $C(O_2 \setminus O) \subseteq O_2 \setminus O$, and hence $C(O_2 \setminus O) \subseteq (O_2 \setminus O) \setminus O_1$. Then indeed $O_1 \cap C(O_2 \setminus O) = \emptyset$, because of Lemma 35. For sufficiency, consider any O, O_1, O_2 in \mathcal{Q} such that $O \subseteq O_2 \setminus O_1$ and $C(O_2) \subseteq O_1$, implying that $(O_2 \setminus O_1) \cap C(O_2) = \emptyset$. Let $\tilde{O} := O$, $\tilde{O}_1 := O_2 \setminus O_1$ and $\tilde{O}_2 := O_2$, then $\tilde{O} \subseteq \tilde{O}_1$ and $\tilde{O}_1 \cap C(\tilde{O}_2) = \emptyset$. Due to Equation (4), then $\tilde{O}_1 \cap C(\tilde{O}_2 \setminus \tilde{O}) = \emptyset$, or equivalently, $C(\tilde{O}_2 \setminus \tilde{O}) \subseteq (\tilde{O}_2 \setminus \tilde{O}) \setminus \tilde{O}_1$, because $C(\tilde{O}_2 \setminus \tilde{O}) \subseteq \tilde{O}_2 \setminus \tilde{O}$. Then $C(O_2 \setminus O) \subseteq (O_2 \setminus O) \setminus (O_2 \setminus O_1)$. But $(O_2 \setminus O) \setminus (O_2 \setminus O_1) = O_2 \cap O_1$, implying that indeed $C(O_2 \setminus O) \subseteq O_2 \cap O_1 \subseteq O_1$. \square

Proof of Proposition 3. For (i), we consider any O in \mathcal{Q} , λ in $\mathbb{R}_{>0}$ and u in \mathcal{V} , and prove that $\lambda C(O) = C(\lambda O)$ and $C(O) + \{u\} = C(O + \{u\})$. That $\lambda C(O) \subseteq C(\lambda O)$ follows at once from Axiom C_{4a}. For the converse inclusion, let $O' := \lambda O$, then it follows from Axiom C_{4a} that $\frac{1}{\lambda} C(O') \subseteq C(\frac{1}{\lambda} O')$, which is equivalent to $\frac{1}{\lambda} C(\lambda O) \subseteq C(O)$. Next, that $C(O) + \{u\} \subseteq C(O + \{u\})$ follows directly from Axiom C_{4b}. For the converse inclusion, let

$O' := O + \{u\}$, then it follows from Axiom C_{4b} that $C(O') + \{-u\} \subseteq C(O' + \{-u\})$, which is equivalent to $C(O + \{u\}) \subseteq C(O) + \{u\}$.

For (ii), fix any u_1 and u_2 in \mathcal{V} such that $u_1 \leq u_2$, any O in \mathcal{Q} and any v in $O \setminus \{u_1, u_2\}$. The proof is trivial if $u_1 = u_2$, so assume that $u_1 < u_2$, whence $u_1 \notin C(\{u_1, u_2\})$ by Axiom C₂. To prove (ii)a, assume that $u_2 \in O$ and $v \notin C(O \cup \{u_1\})$, then by version (3) [with $\tilde{O}_1 := \{u_1\}$, $\tilde{O}_2 := \{u_1, u_2\}$ and $\tilde{O} := O \cup \{u_1\}$] of Axiom C_{3a}, we conclude that $u_1 \notin C(O \cup \{u_1\})$ and therefore $\{u_1, v\} \cap C(O \cup \{u_1\}) = \emptyset$. Version (4) [with $\tilde{O}_1 := \{u_1, v\}$, $\tilde{O} := \{u_1\}$ and $\tilde{O}_2 := O \cup \{u_1\}$] of Axiom C_{3b} then implies that $v \notin C(O \setminus \{u_1\})$. Version (3) [with $\tilde{O}_1 := \{v\}$, $\tilde{O}_2 := O \setminus \{u_1\}$ and $\tilde{O} := O$] of Axiom C_{3a} then implies that indeed $v \notin C(O)$. For (ii)b, assume that $u_1 \in O$ and $v \notin C(O)$, then by version (3) [with $\tilde{O}_1 := \{u_1\}$ for the first assumption and $\tilde{O}_1 := \{v\}$ for the second one, $\tilde{O}_2 := O$ and $\tilde{O} := \{u_2\} \cup O$] of Axiom C_{3a}, the first assumption implies that $u_1 \notin C(\{u_2\} \cup O)$, and the second assumption implies similarly that $v \notin C(\{u_2\} \cup O)$, whence $\{u_1, v\} \cap C(\{u_2\} \cup O) = \emptyset$. By version (4) [with $\tilde{O}_1 := \{u_1, v\}$, $\tilde{O}_2 := \{u_2\} \cup O$ and $\tilde{O} := \{u_1\}$] of Axiom C_{3b}, this implies that indeed $v \notin C(\{u_2\} \cup O \setminus \{u_1\})$.

For (iii), consider any O, O' in \mathcal{Q} such that $C(O) \subseteq O' \subseteq O$. Let $\tilde{O}_1 := C(O)$, $\tilde{O}_2 := O$ and $\tilde{O} := O \setminus O'$, and apply Axiom C_{3b} [we have $C(\tilde{O}_2) = C(O) = \tilde{O}_1$, so $C(\tilde{O}_2) \subseteq \tilde{O}_1$, and also $\tilde{O} = \tilde{O} \setminus O' \subseteq \tilde{O}_2 \setminus C(O) = \tilde{O}_2 \setminus \tilde{O}_1$] to find that $C(\tilde{O}_2 \setminus \tilde{O}) \subseteq \tilde{O}_1$, or equivalently, $C(O') \subseteq C(O)$. Conversely, consider any $u \in C(O)$ and assume *ex absurdo* that $u \notin C(O')$, then we infer from version (3) [with $\tilde{O}_1 := \{u\}$, $\tilde{O}_2 := O'$ and $\tilde{O} := O$] of Axiom C_{3a} that $u \notin C(O)$, a contradiction.

(iv) is an immediate consequence of (iii) [with $O' = C(O)$]. \square

Proof of Proposition 4. For (i), irreflexivity is Axiom S₁. For transitivity, consider any O_1, O_2 and O_3 in \mathcal{Q} , and assume that $O_1 \triangleleft O_2$ and $O_2 \triangleleft O_3$. Use Axiom S_{3a} to infer that then $O_1 \triangleleft O_1 \cup O_2 \cup O_3$ and $O_2 \triangleleft O_1 \cup O_2 \cup O_3$. By Definition 2(ii) this is equivalent to $O_1 \cup O_2 \triangleleft O_1 \cup O_2 \cup O_3$. Now, use Axiom S_{3b} [with $\tilde{O}_1 := O_1 \cup O_2$, $\tilde{O}_2 := O_1 \cup O_2 \cup O_3$ and $\tilde{O} := O_2 \setminus (O_1 \cup O_3) \subseteq \tilde{O}_1$; then $\tilde{O}_1 \setminus \tilde{O} = (O_1 \cup O_2) \setminus (O_2 \setminus (O_1 \cup O_3)) = O_1 \cup (O_2 \cap O_3)$ and $\tilde{O}_2 \setminus \tilde{O} = (O_1 \cup O_2 \cup O_3) \setminus (O_2 \setminus (O_1 \cup O_3)) = O_1 \cup O_3$] to infer that $O_1 \cup (O_2 \cap O_3) \triangleleft O_1 \cup O_3$, and therefore $O_1 \triangleleft O_1 \cup O_3$ using Definition 2(ii). Then indeed $O_1 \triangleleft O_3$ by Definition 2(i).

For (ii), consider any O_1 and O_2 in \mathcal{Q} , any u in \mathcal{V} and any λ in $\mathbb{R}_{>0}$. For the direct implication, assume that $O_1 \triangleleft O_2$. Then $\lambda O_1 \triangleleft \lambda O_2$ by Axiom S_{4a} and therefore indeed $\lambda O_1 + \{u\} \triangleleft \lambda O_2 + \{u\}$ by Axiom S_{4b}. Conversely, for the converse implication, assume that $\lambda O_1 + \{u\} \triangleleft \lambda O_2 + \{u\}$. Then $\lambda O_1 \triangleleft \lambda O_2$ by Axiom S_{4b} and therefore indeed $O_1 \triangleleft O_2$ by Axiom S_{4a}. \square

Proof of Proposition 5. We first prove the first statement. We show that \triangleleft_C satisfies the Axioms S₁–S₄.

For Axiom S₁, consider any O in \mathcal{Q} . Since $C(O) \neq \emptyset$ by Axiom C₁, we infer that $C(O) \not\subseteq \emptyset = O \setminus O$, whence indeed $O \not\triangleleft_C O$.

For Axiom S₂, assume that $u < v$. Then $\{v\} = C(\{u, v\})$ by Axiom C₂, so $C(\{u, v\}) \subseteq \{v\} \setminus \{u\}$. Then indeed $\{u\} \triangleleft_C \{v\}$.

For Axiom S_{3a}, assume that $O_1 \triangleleft_C O_2$ and $O_2 \subseteq O$. Then $O_1 \cap C(O_1 \cup O_2) = \emptyset$ by Definition 3. Use version (3) of Axiom C_{3a} [with $\tilde{O}_1 := O_1$, $\tilde{O}_2 := O_1 \cup O_2$ and $\tilde{O} := O_1 \cup O$; then $\tilde{O}_1 \subseteq \tilde{O}_2 \subseteq \tilde{O}$ since $O_2 \subseteq O$] to infer that then $O_1 \cap C(O_1 \cup O) = \emptyset$, whence indeed $O_1 \triangleleft_C O$.

For Axiom S_{3b}, assume that $O_1 \triangleleft_C O_2$ and $O \subset O_1$. Then $O_1 \cap C(O_1 \cup O_2) = \emptyset$ by Definition 3. Use version (4) of Axiom C_{3b} [with $\tilde{O}_1 := O_1$, $\tilde{O}_2 := O_1 \cup O_2$ and $\tilde{O} := O$; then $\tilde{O} \subseteq \tilde{O}_1$] to infer that $O_1 \cap C((O_1 \cup O_2) \setminus O) = \emptyset$, and hence in particular that $(O_1 \setminus O) \cap C((O_1 \cup O_2) \setminus O) = \emptyset$. Then $O_1 \setminus O \triangleleft_C (O_1 \setminus O) \cup (O_2 \setminus O)$. Use Definition 2(i) [with

$\tilde{O}_1 := O_1 \setminus O$, $\tilde{O}_2 := (O_1 \setminus O) \cup (O_2 \setminus O)$ and $\tilde{O} := O_2 \setminus O$; then $\tilde{O}_2 \setminus \tilde{O}_1 \subseteq \tilde{O} \subseteq \tilde{O}_1 \cup \tilde{O}_2$] to infer that then indeed $O_1 \setminus O \triangleleft_C O_2 \setminus O$.

For Axioms S_{4a} and S_{4b}, assume that $O_1 \triangleleft_C O_2$. We need to show that then $\lambda O_1 + \{u\} \triangleleft_C \lambda O_2 + \{u\}$. To do so, from $O_1 \triangleleft_C O_2$ infer that $O_1 \cap C(O_1 \cup O_2) = \emptyset$, so $(\lambda O_1 + \{u\}) \cap (\lambda C(O_1 \cup O_2) + \{u\}) = \emptyset$. Then use Proposition 3(i) to infer that $(\lambda O_1 + \{u\}) \cap C(\lambda(O_1 \cup O_2) + \{u\}) = \emptyset$, whence indeed $\lambda O_1 + \{u\} \triangleleft_C \lambda O_2 + \{u\}$.

For the second statement, we need to show that C_{\triangleleft} satisfies the Axioms C₁–C₄.

For Axiom C₁, consider any O in \mathcal{Q} . Then $O \not\triangleleft O$ by Axiom S₁, so $\{u\} \not\triangleleft O$ for some u in O , by Definition 2(ii). Then $u \in C_{\triangleleft}(O) := O \setminus \bigcup\{O' \subseteq O : O' \triangleleft O\}$, so $C_{\triangleleft}(O)$ is indeed non-empty.

For Axiom C₂, assume that $u < v$. Then $\{u\} \triangleleft \{v\}$ by Axiom S₂, so $C_{\triangleleft}(\{u, v\}) \cap \{u\} = \emptyset$. Since $C_{\triangleleft}(\{u, v\})$ is non-empty, as shown above, then necessarily $v \in C_{\triangleleft}(\{u, v\})$, so indeed $\{v\} = C_{\triangleleft}(\{u, v\})$.

For version (3) of Axiom C_{3a}, assume that $O_1 \cap C_{\triangleleft}(O_2) = \emptyset$ and $O_1 \subseteq O_2 \subseteq O$. Since $O_1 \subseteq O_2$, then $O_2 = O_1 \cup O_2$, so $O_1 \cap C_{\triangleleft}(O_1 \cup O_2) = \emptyset$, whence $O_1 \triangleleft O_2$. Then $O_1 \triangleleft O$ by Axiom S_{3a}, so $O_1 \cap C_{\triangleleft}(O \cup O_1) = \emptyset$, and since $O \cup O_1 = O$, indeed $O_1 \cap C_{\triangleleft}(O) = \emptyset$.

For version (4) of Axiom C_{3b}, assume that $O_1 \cap C_{\triangleleft}(O_2) = \emptyset$ and $O \subseteq O_1$. We need to prove that then $O_1 \cap C_{\triangleleft}(O_2 \setminus O) = \emptyset$. From $O_1 \cap C_{\triangleleft}(O_2) = \emptyset$, and using that $C_{\triangleleft}(O_2) = O_2 \setminus \bigcup\{O' \subseteq O_2 : O' \triangleleft O_2\}$, infer that $O_1 \cap O_2 \subseteq \bigcup\{O' \subseteq O_2 : O' \triangleleft O_2\}$, whence $O_1 \cap O_2 \triangleleft O_2$, taking into account Definition 2(ii). Call $O' := O \cap O_2$, then $(O_1 \cap O_2) \setminus O' = (O_1 \cap O_2) \setminus O$ and $O_2 \setminus O' = O_2 \setminus O$. Also $O' \subseteq O_1 \cap O_2$, so either $O' = O_1 \cap O_2$ or $O' \subset O_1 \cap O_2$. If $O' = O_1 \cap O_2$ then $(O_1 \cap O_2) \setminus O' = \emptyset$ whence $(O_1 \cap O_2) \setminus O' \cap C(O_2 \setminus O') = \emptyset$ trivially. But since $C(O_2 \setminus O') \subseteq O_2 \setminus O'$, then also $O_1 \cap C(O_2 \setminus O') = \emptyset$ and hence indeed $O_1 \cap C(O_2 \setminus O) = \emptyset$. So assume $O' \subset O_1 \cap O_2$. Then Axiom S_{3b} implies that $(O_1 \cap O_2) \setminus O' \triangleleft O_2 \setminus O'$, or equivalently, replacing O' with $O \cap O_2$, that $(O_1 \cap O_2) \setminus (O \cap O_2) \triangleleft O_2 \setminus (O \cap O_2)$. Infer that $(O_1 \cap O_2) \setminus (O \cap O_2) = (O_1 \cap O_2) \setminus O$ and $O_2 \setminus (O \cap O_2) = O_2 \setminus O$, and therefore $(O_1 \cap O_2) \setminus O \triangleleft O_2 \setminus O$. Then $(O_1 \cap O_2) \setminus O \cap C_{\triangleleft}(O_2 \setminus O) = \emptyset$, and since $C_{\triangleleft}(O_2 \setminus O) \subseteq O_2 \setminus O$, we find that indeed $O_1 \cap C_{\triangleleft}(O_2 \setminus O) = \emptyset$.

For Axioms C_{4a} and C_{4b}, assume that $O_1 \subseteq C_{\triangleleft}(O_2)$. We need to show that then $\lambda O_1 + \{u\} \subseteq C_{\triangleleft}(\lambda O_2 + \{u\})$. To do so, from $O_1 \subseteq C_{\triangleleft}(O_2)$ infer that $O_1 \subseteq O_2$ and hence $O_2 = O_1 \cup O_2$, whence $C_{\triangleleft}(O_1 \cup O_2) \subseteq O_2 \setminus O_1$. Then $O_1 \triangleleft O_2$, so $\lambda O_1 + \{u\} \triangleleft \lambda O_2 + \{u\}$ by Axioms S_{4a} and S_{4b}, whence $C_{\triangleleft}(\lambda(O_1 \cup O_2) + \{u\}) \subseteq \lambda(O_2 \setminus O_1) + \{u\}$. Taking into account that $O_1 \cup O_2 = O_2$, then indeed $\lambda O_1 + \{u\} \subseteq C_{\triangleleft}(\lambda O_2 + \{u\})$. \square

Proof of Proposition 7. The map f we are looking for is $f: \mathcal{C} \rightarrow \mathcal{S}: C \mapsto f(C) := \triangleleft_C$. Due to Proposition 2, f is an invertible function with inverse $f^{-1}: \mathcal{S} \rightarrow \mathcal{C}: \triangleleft \mapsto f^{-1}(\triangleleft) := C_{\triangleleft}$, and hence f is a bijection between \mathcal{C} and \mathcal{S} .

For the first statement, we prove that $(\mathcal{C}; \subseteq)$ and $(\mathcal{S}; \subseteq_{\mathcal{S}})$ are order isomorphic. We first show that $C_1 \subseteq C_2 \Rightarrow f(C_1) \subseteq_{\mathcal{S}} f(C_2)$. So consider any two choice functions C_1 and C_2 such that $C_1 \subseteq C_2$. Then $C_1(O) \supseteq C_2(O)$ for all O in \mathcal{Q} . Consider any O_1 and O_2 in \mathcal{Q} and assume that $O_1 \triangleleft_{C_1} O_2$. Then $O_1 \cap C_1(O_1 \cup O_2) = \emptyset$, and hence also $O_1 \cap C_2(O_1 \cup O_2) = \emptyset$. But then $O_1 \triangleleft_{C_2} O_2$, so indeed $\triangleleft_{C_1} \subseteq_{\mathcal{S}} \triangleleft_{C_2}$. To show that $f(C_1) \subseteq_{\mathcal{S}} f(C_2) \Rightarrow C_1 \subseteq C_2$, it suffices to consider any two choice relation \triangleleft_1 and \triangleleft_2 such that $\triangleleft_1 \subseteq_{\mathcal{S}} \triangleleft_2$, and prove that $C_{\triangleleft_1} \subseteq C_{\triangleleft_2}$, because the map f is a bijection. So assume that $\triangleleft_1 \subseteq_{\mathcal{S}} \triangleleft_2$. Then $O_1 \triangleleft_1 O_2 \Rightarrow O_1 \triangleleft_2 O_2$ for all O_1 and O_2 in \mathcal{Q} . Consider any O in \mathcal{Q} , and any u in O such that $u \notin C_{\triangleleft_1}(O) = O \setminus \bigcup\{O' \subseteq O : O' \triangleleft_1 O\}$. Then $u \in \bigcup\{O' \subseteq O : O' \triangleleft_1 O\}$, but since $\triangleleft_1 \subseteq_{\mathcal{S}} \triangleleft_2$, then also $u \in \bigcup\{O' \subseteq O : O' \triangleleft_2 O\}$. Then $u \notin C_{\triangleleft_2}(O)$, so indeed $C_{\triangleleft_1} \subseteq C_{\triangleleft_2}$.

For the second statement, Proposition 5 guarantees that the same map is a bijection between $\bar{\mathcal{C}}$ and $\bar{\mathcal{D}}$. Since we have already shown that both f and its inverse f^{-1} preserve the order, this completes the proof. \square

Proof of Proposition 8. We show that $\inf \mathcal{C}'$ satisfies the rationality axioms of Definition 4:

- C₁ Consider any C in \mathcal{C}' [always possible since $\mathcal{C}' \neq \emptyset$] and any O in \mathcal{Q} , then $\emptyset \subset C(O) \subseteq \inf \mathcal{C}'(O)$.
- C₂ Note that, for all C in \mathcal{C}' , $C(\{u, v\}) = \{v\}$ whenever $u < v$, implying that $\inf \mathcal{C}'(\{u, v\}) = \bigcup_{C \in \mathcal{C}'} C(\{u, v\}) = \{v\}$.
- C_{3a} Fix any O, O_1, O_2 in \mathcal{Q} such that $\bigcup_{C \in \mathcal{C}'} C(O_2) = \inf \mathcal{C}'(O_2) \subseteq O_2 \setminus O_1$ and $O_1 \subseteq O_2 \subseteq O$. This implies that for all C in \mathcal{C}' , $C(O_2) \subseteq O_2 \setminus O_1$ and therefore also $C(O) \subseteq O \setminus O_1$, because C is coherent. So indeed $\inf \mathcal{C}'(O) = \bigcup_{C \in \mathcal{C}'} C(O) \subseteq O \setminus O_1$.
- C_{3b} Fix any O, O_1, O_2 in \mathcal{Q} such that $\inf \mathcal{C}'(O_2) = \bigcup_{C \in \mathcal{C}'} C(O_2) \subseteq O_1$ and $O \subseteq O_2 \setminus O_1$. This means that for all C in \mathcal{C}' , $C(O_2) \subseteq O_1$, and therefore also $C(O_2 \setminus O) \subseteq O_1$, because C is coherent. So indeed $\inf \mathcal{C}'(O_2 \setminus O) = \bigcup_{C \in \mathcal{C}'} C(O_2 \setminus O) \subseteq O_1$.
- C₄ Consider any C in \mathcal{C}' , λ in $\mathbb{R}_{>0}$, u in \mathcal{V} and O in \mathcal{Q} , and derive from Proposition 3(i) that $\lambda C(O) + \{u\} = C(\lambda O + \{u\})$. Hence $\lambda \inf \mathcal{C}'(O) + \{u\} = \lambda \bigcup_{C \in \mathcal{C}'} C(O) + \{u\} = \bigcup_{C \in \mathcal{C}'} (\lambda C(O) + \{u\}) = \bigcup_{C \in \mathcal{C}'} C(\lambda O + \{u\}) = \inf \mathcal{C}'(\lambda O + \{u\})$, which implies that $\inf \mathcal{C}'$ satisfies axioms C_{4a} and C_{4b}. \square

Proof of Proposition 9. Consider any set of coherent sets of desirable options $\mathcal{D}' \subseteq \bar{\mathcal{D}}$, and its infimum $\inf \mathcal{D}' = \bigcap \mathcal{D}'$, which is of course also a set of desirable options. We show that $\inf \mathcal{D}'$ is coherent, meaning that it satisfies the rationality axioms of Definition 7:

- D₁ Since $0 \notin D$ for all D in \mathcal{D}' , also $0 \notin \inf \mathcal{D}'$.
- D₂ Fix any u in $\mathcal{V}_{>0}$, so $u \in D$ for all D in \mathcal{D}' , implying that indeed $u \in \inf \mathcal{D}'$.
- D₃ Consider any λ in $\mathbb{R}_{>0}$ and any $u \in \inf \mathcal{D}'$, meaning that $u \in D$ for all D in \mathcal{D}' . Then also $\lambda u \in D$ for all D in \mathcal{D}' , implying that indeed $\lambda u \in \inf \mathcal{D}'$.
- D₄ Consider any $u, v \in \inf \mathcal{D}'$, meaning that u and v belong to all D in \mathcal{D}' . Hence $u + v \in D$ for all D in \mathcal{D}' , implying that indeed $u + v \in \inf \mathcal{D}'$.

Now, as a convex cone, $\mathcal{V}_{>0}$ satisfies D₃ and D₄, and by definition it satisfies D₁ and D₂. So $\mathcal{V}_{>0}$ is coherent, and by Axiom D₂ it is included in any other coherent set of desirable options. \square

Proof of Proposition 10. Let, for ease of notation, $D' := \text{posi}(D \cup \{-u\})$. It is clear that D' satisfies Axioms D₂–D₄, so we only need to prove that $0 \notin D'$. Assume *ex absurdo* that $0 \in D'$. Since $0 \notin D$ and $0 \neq u$, there must be v in D and λ in $\mathbb{R}_{>0}$ such that $v + \lambda(-u) = 0$, implying that $u = \frac{1}{\lambda}v \in D$ [Axiom D₃], a contradiction.

Next, consider any maximal coherent set of desirable options D and fix any option u in $\mathcal{V} \setminus \{0\}$ such that $-u \notin D$. Assume *ex absurdo* that also $u \notin D$, then $D' := \text{posi}(D \cup \{-u\}) \supset D$ is a coherent set of desirable options by the first part, contradicting the maximality of D .

Conversely, consider any coherent set of desirable options D that satisfies Equation (5), and any coherent set of desirable options $D' \supseteq D$. Consider any u in D' then clearly $u \neq 0$ [Axiom D₁], so we infer from (5) that $u \in D$ or $-u \in D$. Assume *ex absurdo* that $-u \in D$, then also $-u \in D'$, which together with $u \in D'$ implies that $0 = u + (-u) \in D'$ [Axiom D₄], which contradicts the coherence of D' [Axiom D₁]. Hence $u \in D$, which implies that $D' = D$, so D is indeed maximal. \square

Proof of Proposition 11. We have to prove that the set $\{\hat{D} \in \bar{\mathcal{D}} : D \subseteq \hat{D}\}$ has a maximal element. This will follow directly from Zorn's Lemma if we can show that any chain \mathbb{K} in

this poset has a greatest element. It is a matter of straightforward verification of the axioms to see that $\bigcup \mathbb{K}$ is a coherent set of desirable options, so $\bigcup \mathbb{K}$ is this greatest element.

Let us now establish the dual atomicity. By definition, $D \subseteq \bigcap \hat{D}_D$, so we concentrate on proving that $\bigcap \hat{D}_D \subseteq D$. Consider any u in $\bigcap \hat{D}_D$, meaning that $u \in \hat{D}$ for all $\hat{D} \in \hat{D}_D$. Assume *ex absurdo* that $u \notin D$, then $-u \in \text{posi}(D \cup \{-u\})$, and $\text{posi}(D \cup \{-u\})$ is a coherent set of desirable options by Proposition 10. Consider any maximal coherent set of desirable options \hat{D} that dominates this set [there is such a coherent maximal set by Proposition 11], then also $-u \in \hat{D}$ and therefore $u \notin \hat{D}$ [use Axioms D_4 and D_1]. But since the maximal \hat{D} also dominates D , this is a contradiction. \square

Proof of Proposition 12. We first show that D_C is coherent:

- D_1 Assume *ex absurdo* that $0 \in D_C$, meaning that $0 \notin C(\{0\})$, contradicting the coherence of C [Axiom C_1].
- D_2 Consider any option u in $\mathcal{V}_{>0}$, so $0 < u$. The coherence of C [Axiom C_2] then implies that $\{u\} = C(\{0, u\})$, whence $0 \notin C(\{0, u\})$, implying that indeed $u \in D_C$.
- D_3 Consider any λ in $\mathbb{R}_{>0}$, and any u in D_C , meaning that $0 \notin C(\{0, u\})$. Proposition 3(i) with $v = 0$ then implies that $0 = \lambda 0 \notin C(\{0, \lambda u\})$, whence indeed $\lambda u \in D_C$.
- D_4 Consider any u, v in D , meaning that $0 \notin C(\{0, u\})$ and $0 \notin C(\{0, v\})$. Proposition 3(i) with $\mu = 1$ then implies that $u \notin C(\{u, u+v\})$ and $v \notin C(\{v, u+v\})$, and consequently that $\{0, u, v\} \cap C(\{0, u, v, u+v\}) = \emptyset$ [use version (3) of Axiom C_3a three times with $O := \{0, u, v, u+v\}$: once with $O_1 := \{0\}$ and $O_2 := \{0, u\}$, once with $O_1 := \{u\}$ and $O_2 := \{u, u+v\}$, and once with $O_1 := \{v\}$ and $O_2 := \{v, u+v\}$]. This implies that $0 \notin C(\{0, u, v\})$ by version (4) [with $O_3 := \{0, u, v\}$, $O_2 := \{0, u, v, u+v\}$ and $O := \{u, v\}$] of Axiom C_3b , whence indeed $u+v \in D_C$.

We complete the proof by showing that C and D_C are compatible: $v \notin C(\{u, v\}) \Leftrightarrow u-v \in D_C$ for all u, v in \mathcal{V} . For the direct implication, consider any u, v in \mathcal{V} such that $v \notin C(\{u, v\})$. Proposition 3(i) with $\lambda = 1$ then guarantees that $0 \notin C(\{0, u-v\})$, implying that indeed $u-v \in D_C$. For the converse implication, consider any u, v in \mathcal{V} such that $u-v \in D_C$. Then $0 \notin C(\{0, u-v\})$, implying that indeed $v \notin C(\{u, v\})$, by Proposition 3(i) with $\lambda = 1$. \square

Proof of Proposition 13. The proof is structured as follows: we show (a) that C_D is compatible with D ; (b) that C_D is coherent; and (c) that $C_D \subseteq C$ for all $C \in \bar{C}_D$.

- (a) First, we show that C_D is compatible with D : Consider any u, v in \mathcal{V} then it follows from the definition of C_D that indeed

$$v \notin C_D(\{u, v\}) \Leftrightarrow (\exists w \in \{u, v\}) w - v \in D \Leftrightarrow (u - v \in D \text{ or } v - v \in D) \Leftrightarrow u - v \in D,$$

where the last equivalence follows from $0 \notin D$, because D is coherent [Axiom D_1].

- (b) Next, we show that C_D is coherent:

- C_1 Consider any O in \mathcal{Q} . Since $|O|$ is finite and \prec is a strict partial order, we know that there is at least one maximal element u_m for \prec , meaning that $(\forall v \in O) u_m \not\prec v$, or equivalently $(\forall v \in O) v - u_m \notin D$. Hence $u_m \in C_D(O)$.
- C_2 Consider any u, v in \mathcal{V} such that $u < v$. Then $0 < v - u$, whence $v - u \in D$ by Axiom D_2 . So $u \notin C_D(\{u, v\})$ by compatibility. Since we already know that $C_D(\{u, v\}) \neq \emptyset$, we must therefore have that indeed $\{v\} = C_D(\{u, v\})$.
- (3) Consider any option sets O, O_1, O_2 in \mathcal{Q} such that $O_1 \subseteq O_2 \subseteq O$ and $O_1 \cap C_D(O_2) = \emptyset$, implying that $u \notin C_D(O_2)$ for all u in O_1 , or equivalently, that $(\forall u \in O_1)(\exists v \in O_2) v - u \in D$. Since $O_2 \subseteq O$, this implies that $(\forall u \in O_1)(\exists v \in O) v - u \in D$, or indeed $O_1 \cap C_D(O) = \emptyset$.
- (4) Consider any option sets O, O_1, O_2 in \mathcal{Q} such that $O \subseteq O_1$ and $O_1 \cap C_D(O_2) = \emptyset$, implying that $(\forall u \in O_1)(\exists v \in O_2) v - u \in D$. We need to prove that $(\forall u \in O_1)(\exists v \in$

$O_2 \setminus O)v - u \in D$. We may assume without loss of generality that $O = \{v_1, \dots, v_n\}$ with n in \mathbb{N}_0 . Consider any u in O_1 , then we know that $u_1 - u \in D$ for some u_1 in O_2 . Then either $u_1 \in O_2 \setminus O$ —in which case the proof is finished—or $u_1 \in O$, so we may assume that $u_1 = v_1$ without loss of generality. But then, since $O \subseteq O_1$, we know that $u_2 - v_1 \in D$ for some u_2 in O_2 . By coherence of D [Axiom D_1], we know that $u_2 \neq v_1$, and therefore also that u_2 in $O_2 \setminus \{v_1\}$. Then either $u_2 \in O_2 \setminus O$ —in which case the proof is finished—or $u_2 \in O \setminus \{v_1\}$, so we may assume that $u_2 = v_2$ without loss of generality. This tells us that $v_2 - v_1 \in D$. But then, again, since $O \subseteq O_1$, we know that $u_3 - v_2 \in D$ for some u_3 in O_2 . The coherence of D [Axiom D_4] then implies that also $u_3 - v_1 = u_3 - v_2 + v_2 - v_1 \in D$, and hence that [use Axiom D_1 twice] $u_3 \notin \{v_1, v_2\}$, so we know that u_3 in $O_2 \setminus \{v_1, v_2\}$. When we proceed in this way, we are guaranteed to find, after $k \leq n + 1$ steps, some u_k in $O_2 \setminus O$ such that $u_k - u \in D$.

C₄ Consider any O_1, O_2 in \mathcal{Q} such that $O_1 \subseteq C_D(O_2)$, meaning that $(\forall u \in O_1)(\forall v \in O_2)v - u \notin D$.

C_{4a} It follows that $(\forall u \in O_1)(\forall v \in O_2)\lambda v - \lambda u \notin D$ for any λ in $\mathbb{R}_{>0}$, whence indeed $\lambda O_1 \subseteq C_D(\lambda O_2)$.

C_{4b} Also $(\forall u \in O_1)(\forall v \in O_2)(v + w) - (u + w) \notin D$ for any w in \mathcal{V} , whence indeed $O_1 + \{w\} \subseteq C_D(O_2 + \{w\})$.

(c) Finally, we show that $C_D \subseteq C$ for all $C \in \bar{C}_D$: Consider any C in \bar{C}_D and O in \mathcal{Q} , then we have to prove that $C(O) \subseteq C_D(O)$. Consider any u in O and assume that $u \notin C_D(O)$. Then $v - u \in D$ for some v in O , implying that $u \notin C(\{u, v\})$ because C is compatible with D . The coherence of C [version (3) of Axiom C_{3a} with $\tilde{O}_1 := \{u\}$, $\tilde{O}_2 := \{u, v\}$ and $\tilde{O} := O$] then implies that also $u \notin C(O)$. \square

Proof of Proposition 14. Because $D_v \in \bar{D}$, we have that C_{D_v} is a coherent choice function by Proposition 13, and it then follows that $C_v \subseteq C_{D_v}$. We show that $C_{D_v} \subseteq C_v$, or in other words, that

$$(\forall O \in \mathcal{Q})(\forall u \in O)u \notin C_{D_v}(O) \Rightarrow u \notin C_v(O)$$

Fix any O in \mathcal{Q} and any u in O such that $u \notin C_{D_v}(O)$, meaning that $u < v$ for some v in O by the definition of C_{D_v} . Then, due to Axiom C₂, we find that $u \notin C(\{u, v\})$, and using version (3) of Axiom C_{3a} [with $\tilde{O}_1 := \{u\}$, $\tilde{O}_2 := \{u, v\}$ and $\tilde{O} := O$] we infer that $u \notin C_v(O)$. \square

Proof of Proposition 15. Recall in advance that $\text{inf}C'$ is a coherent choice function by Proposition 8, and that $\text{inf}\{D_C : C \in C'\}$ is a coherent set of desirable options by Proposition 9.

For the first statement, consider any u in \mathcal{V} , and observe that

$$\begin{aligned} u \in D_{\text{inf}C'} &\Leftrightarrow 0 \notin \text{inf}C'(\{0, u\}) \Leftrightarrow 0 \notin \bigcup_{C \in C'} C(\{0, u\}) \\ &\Leftrightarrow (\forall C \in C') 0 \notin C(\{0, u\}) \\ &\Leftrightarrow (\forall C \in C') u \in D_C \\ &\Leftrightarrow u \in \bigcap \{D_C : C \in C'\} = \text{inf}\{D_C : C \in C'\}, \end{aligned}$$

where the first and fourth equivalences follow from Proposition 12.

For the second statement, fix any O in \mathcal{Q} and u in \mathcal{V} such that $u \in \text{inf}C'(O)$. Then $u \in C(O)$ for some C in C' , from which we infer that $(\exists C \in C')(\forall v \in O)u \in C(\{u, v\})$, by an immediate application of version (3) [use the contraposition with $\tilde{O}_1 := \{u\}$, $\tilde{O}_2 := \{u, v\}$ and $\tilde{O} := O$] of Axiom C_{3a}. By exchanging the quantifiers, we infer as an implication that

$(\forall v \in O)(\exists C \in \mathcal{C}')u \in C(\{u, v\})$. Now recall that

$$\begin{aligned} u \in C(\{u, v\}) &\Leftrightarrow 0 \in C(\{0, v-u\}) && \text{by Axiom C}_{4b} \\ &\Leftrightarrow v-u \notin D_C && \text{by Proposition 12.} \end{aligned}$$

This implies that $(\forall v \in O)v-u \notin \bigcap \{D_C : C \in \mathcal{C}'\}$, which is equivalent to $u \in C_{\inf\{D_C : C \in \mathcal{C}'\}}(O)$, by Proposition 13. The rest of the proof is now immediate. \square

Proof of Proposition 16. Recall in advance that $\inf\{C_D : D \in \mathcal{D}'\}$ is a coherent choice function by Propositions 8 and 13, and that $\inf\mathcal{D}'$ is a coherent set of desirable options by Proposition 9.

For the first statement, consider any u in \mathcal{V} and observe that:

$$\begin{aligned} u \in D_{\inf\{C_D : D \in \mathcal{D}'\}} &\Leftrightarrow 0 \notin (\inf\{C_D : D \in \mathcal{D}'\})(\{0, u\}) = \bigcup_{D \in \mathcal{D}'} C_D(\{0, u\}) && \text{by Pr. 12} \\ &\Leftrightarrow (\forall D \in \mathcal{D}') 0 \notin C_D(\{0, u\}) \\ &\Leftrightarrow (\forall D \in \mathcal{D}') (\exists v \in \{0, u\}) v \in D && \text{by Pr. 13} \\ &\Leftrightarrow (\forall D \in \mathcal{D}') u \in D && \text{by Ax. D}_1 \\ &\Leftrightarrow u \in \inf\mathcal{D}'. \end{aligned}$$

For the second statement, let $\mathcal{C}' := \{C_D : D \in \mathcal{D}'\}$, then we infer from the first statement that $\{D_C : C \in \mathcal{C}'\} = \mathcal{D}'$. Now, use the second statement in Proposition 15. \square

Proof of Corollary 17. It is an immediate consequence of Propositions 16 and 15, by letting $\mathcal{D}' := \{D\}$ and $\mathcal{C}' := \{C\}$. \square

Proof of Proposition 18. Immediate consequence of Propositions 8 and 13. \square

Proof of Proposition 19. Consider any u in \mathcal{V} , then

$$\begin{aligned} u \in D_{C_{\mathcal{D}'}}^M &\Leftrightarrow 0 \notin C_{\mathcal{D}'}^M(\{0, u\}) && \text{by Proposition 12} \\ &\Leftrightarrow (\forall \hat{D} \in \hat{\mathcal{D}}_{\mathcal{D}'}) 0 \notin C_{\hat{D}}(\{0, u\}) && \text{by Definition 11} \\ &\Leftrightarrow (\forall \hat{D} \in \hat{\mathcal{D}}_{\mathcal{D}'}) u \in \hat{D} && \text{by Proposition 12} \\ &\Leftrightarrow u \in \bigcap \hat{\mathcal{D}}_{\mathcal{D}'} \Leftrightarrow u \in D' && \text{by Proposition 11,} \end{aligned}$$

proving the first statement. It then follows from Proposition 12 that D' is compatible with $C_{\mathcal{D}'}^M$, and therefore from Proposition 13 that $C_{D'} = \inf \bar{C}_{\mathcal{D}'} \subseteq C_{\mathcal{D}'}^M$. \square

Proof of Proposition 20. By Proposition 13, it suffices to prove (a) that D_p is coherent; and (b) that $C_p = C_{D_p}$.

For (a), that Axiom D₁ holds, follows from $0 \notin \mathcal{V}_{>0}$ and $E_p(0) = 0$. Axiom D₂ holds by definition. For Axiom D₃, consider any u in D_p and real λ in $\mathbb{R}_{>0}$, then $0 < u$ and therefore $0 < \lambda u$, or $E_p(u) > 0$ and therefore $E_p(\lambda u) = \lambda E_p(u) > 0$, whence indeed $\lambda u \in D_p$. For Axiom D₄, consider any u, v in D_p , then there are three possibilities. The first is that both u and v belong to $\mathcal{V}_{>0}$, and therefore also $u+v \in \mathcal{V}_{>0}$. The second is that both $E_p(u) > 0$ and $E_p(v) > 0$, and therefore also $E_p(u+v) = E_p(u) + E_p(v) > 0$. And the third is that, without loss of generality, $E_p(u) > 0$ and $v \in \mathcal{V}_{>0}$, whence $E_p(v) \geq 0$ and therefore $E_p(u+v) = E_p(u) + E_p(v) > 0$. In all cases therefore indeed $u+v \in D_p$.

For (b), consider any O in \mathcal{Q} , use Equation (8) to find that $C_{D_p}(O) = \{u \in O : (\forall v \in O)v-u \notin D_p\}$ and Equation (10) to find that $v-u \notin D_p \Leftrightarrow (v-u \notin \mathcal{V}_{>0} \text{ and } E_p(v-u) \leq 0)$. It then follows from Equation (11) that indeed $C_p(O) = C_{D_p}(O)$. \square

Proof of Proposition 21. We infer from Definition 12 and Proposition 20 that

$$C_K^E(O) = \bigcup_{p \in K} C_{D_p}(O) = \bigcup_{p \in K} \{u \in O : (\forall v \in O)v - u \notin D_p\},$$

where the last equality follows from Proposition 13. Now

$$\begin{aligned} (\forall v \in O)v - u \notin D_p &\Leftrightarrow (\forall v \in O)(u \not\prec v \text{ and } E_p(v - u) \leq 0) \\ &\Leftrightarrow u \in C_v(O) \text{ and } (\forall v \in O)E_p(v - u) \leq 0, \end{aligned}$$

where the first equivalence follows from Equation (10), and the second from Proposition 14. Hence indeed

$$\begin{aligned} C_K^E(O) &= \bigcup_{p \in K} \{u \in O : (\forall v \in O)E_p(v - u) \leq 0\} \cap C_v(O) \\ &= \{u \in O : (\exists p \in K)(\forall v \in O)E_p(v) \leq E_p(u)\} \cap C_v(O). \quad \square \end{aligned}$$

Proof of Proposition 22. We consider any p in K and prove that $C_{D_p} \subseteq C_{\hat{D}_{\{p\}}}^M$. The proof then follows by taking the infimum $\inf_{p \in K}$ over $p \in K$ on both sides of this inequality. Consider any O in \mathcal{Q} and assume that $u \in C_{\hat{D}_{\{p\}}}^M(O)$, so there is some \hat{D} in $\hat{\mathcal{D}}_{\{p\}}$ such that $u \in C_{\hat{D}}(O)$, or equivalently, $(\forall v \in O)v - u \notin \hat{D}$. Hence $(\forall v \in O)v - u \notin D_p$, because $D_p \subseteq \hat{D}$, and therefore indeed $u \in C_{D_p}(O)$. \square

Proof of Proposition 23. To see that \leq is reflexive, use $w = 0$ in Equation (13). For antisymmetry, consider any $[u], [v]$ in \mathcal{V}/I such that $[u] \leq [v]$ and $[v] \leq [u]$, and therefore $u \leq v + w$ and $v \leq u + w'$ for some $w, w' \in I$. This implies, by repeatedly applying Equation (1) and the transitivity of the vector ordering on \mathcal{V} that $w + w' \geq 0$. Now, by Axiom I₄, $w + w' \in I$, and therefore, by Axiom I₂, $w + w' = 0$, implying that $u \leq v + w$ and $v \leq u - w$, and therefore also $v + w \leq u$, for some $w \in I$. Hence indeed $u = v + w$ and therefore $[u] = [v]$, by the antisymmetry of the vector ordering on \mathcal{V} . For transitivity, consider any $[u], [v], [w]$ in \mathcal{V}/I such that $[u] \leq [v]$ and $[v] \leq [w]$, and therefore $u \leq v + v'$ and $v \leq w + w'$ for some $v', w' \in I$. Then $v + v' \leq w + w' + v'$ by Equation (1), whence also $u \leq w + w' + v'$ by transitivity of the vector ordering on \mathcal{V} . By coherence [Axiom I₄] also $v' + w' \in I$, whence indeed $[u] \leq [w]$ by Equation (13).

Next, we prove that \leq satisfies Eqs. (1) and (2). Consider any $[u_1], [u_2]$ in \mathcal{V}/I such that $[u_1] \leq [u_2]$, implying that $u_1 \leq u_2 + u'_2$ for some $u'_2 \in I$. For (1), consider any v in \mathcal{V} , then also $u_1 + v \leq u_2 + v + u'_2$, implying that indeed $[u_1] + [v] \leq [u_2] + [v]$. The converse implication follows analogously, by adding $-v$ rather than v . For (2), consider any λ in $\mathbb{R}_{>0}$, then also $\lambda u_1 \leq \lambda u_2 + \lambda u'_2$, implying that indeed $\lambda [u_1] \leq \lambda [u_2]$ because $\lambda u'_2 \in I$ by Axiom I₃. The converse implication follows analogously, by multiplying with $\frac{1}{\lambda}$ rather than λ .

We turn to the second statement. For necessity, consider u, v in \mathcal{V} such that $[u] < [v]$, so $[u] \neq [v]$ and $u \leq v + w$ for some w in I . $[u] \neq [v]$ implies that $u \neq v + w$ for all w in I . Taken together with $u \leq v + w$ for some w , this implies that indeed $u < v + w$ for some w . For sufficiency, consider u, v in \mathcal{V} such that $u < v + w$ for some w in I . Assume *ex absurdo* that $[u] \not\prec [v]$, meaning that $[u] = [v]$ or $u \not\leq v + w$ for all w in I . The latter possibility is incompatible with the assumption that $u < v + w$ for some w in I , leaving only the first possibility, which is equivalent to $u = v + w'$ for some w' in I . Then $v + w' < v + w$, implying that $0 < w - w'$ by Equation (1). This contradicts the coherence of I [Axiom I₂] because $w - w'$ is an element of I by Axiom I₄. \square

Proof of Proposition 24. For necessity, observe that

$$D + I = \{u + v : u \in D \text{ and } v \in I\} = \bigcup \{\{u\} + I : u \in D\} = \bigcup \{[u] : u \in D\} = \bigcup D/I,$$

and that compatibility with I guarantees that $D + I = D$. For sufficiency, assume that there is some set of desirable options $D' \subseteq \mathcal{V}/I$ such that $D = \{u : [u] \in D'\}$. Consider any u in D and any v in $I = [0]$, then $u + v \in [u] + [0] = [u + 0] = [u]$, and therefore indeed $u + v \in D$. Then also $D/I = \{[u] : u \in D\} = \{[u] : [u] \in D'\} = D'$. \square

Proof of Proposition 25. For the direct implication, assume that $D \subseteq \mathcal{V}$ is coherent. We show that $D/I = \{[u] : u \in D\} \subseteq \mathcal{V}/I$ satisfies the Axioms D₁–D₄. For Axiom D₁, assume *ex absurdo* that $[0] \in D/I$, implying that $u \in D$ for some u in $[0] = I$, but then $u \in D \cap I$ and hence $D \cap I$ is non-empty, a contradiction. For Axiom D₂, consider any $[u]$ in \mathcal{V}/I such that $[0] < [u]$, meaning that $0 < u + w$ for some w in I . Then $u + w \in D$ for some w in I , by coherence of D [Axiom D₂], implying that $[u] = [u] + I = [u] + [w] = [u + w] \in D/I$. For Axiom D₃, assume that $[u] \in D/I$, and consider any λ in $\mathbb{R}_{>0}$. Then $u \in D$, implying that $\lambda u \in D$ by coherence of D [Axiom D₃], whence $\lambda[u] = [\lambda u] \in D/I$. For Axiom D₄, assume that $[u], [v] \in D/I$, implying that $u, v \in D$. Then $u + v \in D$ by coherence of D [Axiom D₄], whence $[u] + [v] = [u + v] \in D/I$.

For the converse implication, assume that $D/I \subseteq \mathcal{V}/I$ is coherent. We show that $D = \{u : [u] \in D/I\} \subseteq \mathcal{V}$ satisfies the Axioms D₁–D₄. For Axiom D₁, infer that $0 \notin D$ since $[0] \notin D/I$ by coherence of D/I [Axiom D₁]. For Axiom D₂, consider any u in \mathcal{V} such that $0 < u$. Then $[0] < [u]$ by taking for example $w = 0$ in Proposition 23, implying that $[u] \in D/I$ by coherence of D/I [Axiom D₂], whence $u \in D$. For Axiom D₃, assume that $u \in D$, and consider any λ in $\mathbb{R}_{>0}$. Then $[u] \in D/I$, implying that $\lambda[u] = [\lambda u] \in D/I$ by coherence of D/I [Axiom D₃], whence $\lambda u \in D$. For Axiom D₄, assume that $u, v \in D$, implying that $[u], [v] \in D/I$. Then $[u] + [v] \in D/I$ by coherence of D/I [Axiom D₄], implying that $u + v \in D$. \square

Proof of Proposition 26. For the first statement, assume that C is compatible with I , that is, there is some C' on $\mathcal{Q}(\mathcal{V}/I)$ such that $C(O) := \{u \in O : [u] \in C'(O/I)\}$ for all O in $\mathcal{Q}(\mathcal{V})$. Consider any O_1 and O_2 in $\mathcal{Q}(\mathcal{V})$. Infer that

$$\begin{aligned} O_1 \triangleleft_C O_2 &\Leftrightarrow O_1 \cap C(O_1 \cup O_2) = \emptyset \Leftrightarrow O_1 \cap \{u \in O_1 \cup O_2 : [u] \in C'((O_1 \cup O_2)/I)\} = \emptyset \\ &\Leftrightarrow O_1/I \cap C'((O_1 \cup O_2)/I) = \emptyset \\ &\Leftrightarrow O_1/I \triangleleft_{C'} O_2/I. \end{aligned}$$

This shows that the representing choice relation is derived from the representing choice function.

For the second statement, assume that \triangleleft is compatible with I , that is, there is some \triangleleft' on $\mathcal{Q}(\mathcal{V}/I)$ such that $O_1 \triangleleft O_2 \Leftrightarrow O_1/I \triangleleft' O_2/I$ for all O_1 and O_2 in $\mathcal{Q}(\mathcal{V})$. Consider any O in $\mathcal{Q}(\mathcal{V})$. Then $C_{\triangleleft}(O) = O \setminus \bigcup\{O' \subseteq O : O' \triangleleft O\} = O \setminus \bigcup\{O' \subseteq O : O'/I \triangleleft' O/I\}$. Infer that

$$\begin{aligned} u \in C_{\triangleleft}(O) &\Leftrightarrow u \in O \text{ and } u \notin \bigcup\{O' \subseteq O : O'/I \triangleleft' O/I\} \\ &\Leftrightarrow u \in O \text{ and } [u] \notin \bigcup\{\tilde{O} \subseteq O/I : \tilde{O} \triangleleft' O/I\} \\ &\Leftrightarrow u \in O \text{ and } [u] \in O/I \setminus \bigcup\{\tilde{O} \subseteq O/I : \tilde{O} \triangleleft' O/I\} \\ &\Leftrightarrow u \in O \text{ and } [u] \in C_{\triangleleft'}(O/I), \end{aligned}$$

whence C_{\triangleleft} is indeed compatible with I , with representing choice function $C_{\triangleleft'}$. \square

Proof of Proposition 27. Let C' be any representing choice function, and consider any O in $\mathcal{Q}(\mathcal{V})$, then

$$C(O)/I = \{[u] : u \in C(O)\} = \{[u] : u \in O, [u] \in C'(O/I)\} = O/I \cap C'(O/I) = C'(O/I). \quad \square$$

Lemma 36. *Given any two option sets O_1, O_2 in $\mathcal{Q}(\mathcal{V})$, then $O_2/I \setminus O_1/I \subseteq (O_2 \setminus O_1)/I$.*

Proof. Consider any \tilde{u} in $O_2/I \setminus O_1/I$, then $\tilde{u} = [v_2]$ for some v_2 in O_2 and $\tilde{u} \neq [v_1]$ for all v_1 in O_1 , implying that indeed $\tilde{u} = [v_2]$ for some v_2 in $O_2 \setminus O_1$. \square

Lemma 37. *For any choice function C on $\mathcal{Q}(\mathcal{V})$ that is compatible with a coherent set of indifferent options I , and any option set O in $\mathcal{Q}(\mathcal{V})$, we have that $(O \setminus C(O))/I = O/I \setminus C/I(O/I)$.*

Proof. Consider any \tilde{u} in \mathcal{V}/I , then

$$\begin{aligned} \tilde{u} \in (O \setminus C(O))/I &\Leftrightarrow (\exists v \in O \setminus C(O)) \tilde{u} = [v] \\ &\Leftrightarrow (\exists v \in O) (\tilde{u} = [v] \text{ and } v \notin C(O)) \\ &\Leftrightarrow (\exists v \in O) (\tilde{u} = [v] \text{ and } [v] \notin C/I(O/I)) \quad \text{by Proposition 27} \\ &\Leftrightarrow (\exists v \in O) (\tilde{u} = [v] \text{ and } \tilde{u} \notin C/I(O/I)) \\ &\Leftrightarrow \tilde{u} \notin C/I(O/I) \text{ and } (\exists v \in O) \tilde{u} = [v] \\ &\Leftrightarrow \tilde{u} \in O/I \setminus C/I(O/I), \end{aligned}$$

whence indeed $(O \setminus C(O))/I = O/I \setminus C/I(O/I)$. \square

Lemma 38. *For any choice function C on $\mathcal{Q}(\mathcal{V})$ that is compatible with a coherent set of indifferent options I , and any option sets O_1, O_2 in $\mathcal{Q}(\mathcal{V})$ such that $O_1 \subseteq O_2$, we have that $O_1 \cap C(O_2) = \emptyset \Leftrightarrow O_1/I \cap C/I(O_2/I) = \emptyset$.*

Proof. First, assume that $O_1 \cap C(O_2) \neq \emptyset$, and consider any v in $O_1 \cap C(O_2)$. Then $[v] \in O_1/I$ and $[v] \in C(O_2)/I = C/I(O_2/I)$, by Proposition 27. Conversely, assume that $O_1/I \cap C/I(O_2/I) \neq \emptyset$, and consider any \tilde{u} in $O_1/I \cap C/I(O_2/I)$. Then there is some v in O_1 such that $\tilde{u} = [v]$, and we infer from Proposition 27 that $\tilde{u} \in C(O_2)/I$, so also $v \in C(O_2)$. \square

Lemma 39. *For all \tilde{O} in $\mathcal{Q}(\mathcal{V}/I)$, there is some O in $\mathcal{Q}(\mathcal{V})$ such that $O/I = \tilde{O}$.*

Proof. For every element \tilde{u} in the finite set \tilde{O} , we take some vector u in \mathcal{V} such that $[u] = \tilde{u}$, and collect those vectors in an option set O . Then $O \in \mathcal{Q}(\mathcal{V})$ because it is a finite subset of \mathcal{V} , and $O/I = \{[u] : u \in O\} = \tilde{O}$. \square

Proof of Proposition 28. For the direct implication, assume that C is coherent. We show that C/I satisfies Axioms C_1 – C_4b .

For Axiom C_1 , assume *ex absurdo* that $C/I(\tilde{O}) = \emptyset$ for some \tilde{O} in $\mathcal{Q}(\mathcal{V}/I)$. Consider any option set O in $\mathcal{Q}(\mathcal{V})$ such that $O/I = \tilde{O}$, then $C(O) = \{u \in O : [u] \in C/I(\tilde{O})\} = \{u \in O : [u] \in \emptyset\} = \emptyset$, contradicting the coherence of C [Axiom C_1].

For Axiom C_2 , consider $[u], [v]$ in \mathcal{V}/I such that $[u] < [v]$, meaning that $u < v + w$ for some w in I , by Proposition 23. Then $\{v + w\} = C(\{u, v + w\})$ by coherence of C [Axiom C_2], so $\{[v]\} = \{[v + w]\} = C/I(\{[u], [v + w]\}) = C/I(\{[u], [v]\})$, where the first and last equalities hold because w belongs to I .

For version (3) of Axiom C_3a , consider any option sets $\tilde{O}_1, \tilde{O}_2, \tilde{O}$ in $\mathcal{Q}(\mathcal{V}/I)$ such that $\tilde{O}_1 \cap C/I(\tilde{O}_2) = \emptyset$ and $\tilde{O}_1 \subseteq \tilde{O}_2 \subseteq \tilde{O}$. Consider any option set O in $\mathcal{Q}(\mathcal{V})$ such that $O/I = \tilde{O}$, and define the option sets $O_2 := \{u \in O : [u] \in \tilde{O}_2\} \in \mathcal{Q}(\mathcal{V})$ and $O_1 := \{u \in O_2 : [u] \in \tilde{O}_1\} \in \mathcal{Q}(\mathcal{V})$. Then $O_1 \subseteq O_2 \subseteq O$, $O_2/I = \{[v] : v \in O_2\} = \{[v] : v \in O \text{ and } [v] \in \tilde{O}_2\} = O/I \cap \{[v] : [v] \in \tilde{O}_2\} = \tilde{O} \cap \tilde{O}_2 = \tilde{O}_2$ and, similarly, $O_1/I = \{[v] : v \in O_2 \text{ and } [v] \in \tilde{O}_1\} = O_2/I \cap \tilde{O}_1 = \tilde{O}_2 \cap \tilde{O}_1 = \tilde{O}_1$, so we infer from Lemma 38 that $O_1 \cap C(O_2) = \emptyset$, and therefore also that $O_1 \cap C(O) = \emptyset$ by coherence of C [version (3) of Axiom C_3a]. Because $\tilde{O} = O/I$, we conclude from Lemma 38 that indeed $\tilde{O}_1 \cap C/I(\tilde{O}) = \emptyset$.

For version (4) of Axiom C_3b , consider any option sets $\tilde{O}_1, \tilde{O}_2, \tilde{O}$ in $\mathcal{Q}(\mathcal{V}/I)$ such that $\tilde{O}_1 \cap C/I(\tilde{O}_2) = \emptyset$ and $\tilde{O} \subseteq \tilde{O}_1$. Let $\tilde{O}' := \tilde{O}_1 \cap \tilde{O}_2$ and $\tilde{O}'' := \tilde{O}_1 \setminus \tilde{O}_2$, then $\tilde{O} = \tilde{O}' \cup \tilde{O}''$,

and $\tilde{O}' \subseteq \tilde{O}'_1 \subseteq \tilde{O}_2$. So the assumption $\tilde{O}_1 \cap C/I(\tilde{O}_2) = \emptyset$ implies that $\tilde{O}'_1 \cap C/I(\tilde{O}_2) = \emptyset$. Introduce the option sets O', O'_1, O_2 in $\mathcal{Q}(\mathcal{V})$ as follows: O_2 is any option set such that $O_2/I = \tilde{O}_2$, $O'_1 := \{u \in O_2 : [u] \in \tilde{O}'_1\}$ and $O' := \{u \in O'_1 : [u] \in \tilde{O}'\} = \{u \in O_2 : [u] \in \tilde{O}'\}$, then $O' \subseteq O'_1 \subseteq O_2$. Also, $O'_1/I = \tilde{O}'_1$, $O'/I = \tilde{O}'$, and $\tilde{O}'_1 \cap C/I(\tilde{O}_2) = \emptyset$ allows us to infer that $O'_1 \cap C(O_2) = \emptyset$, using Lemma 38. The coherence of C [version (4) of Axiom C₃b] then implies that $O'_1 \cap C(O_2 \setminus O') = \emptyset$, and therefore also $(O'_1 \setminus O') \cap C(O_2 \setminus O') = \emptyset$. This implies that $O'_1 \setminus O' \subseteq (O_2 \setminus O') \setminus C(O_2 \setminus O')$, because $O'_1 \setminus O' \subseteq O_2 \setminus O'$. By Equation (14) and Lemma 37, this implies that $(O'_1 \setminus O')/I \subseteq (O_2 \setminus O')/I \setminus C/I((O_2 \setminus O')/I)$. We now prove that

$$(O'_1 \setminus O')/I = O'_1/I \setminus O'/I$$

and that

$$(O_2 \setminus O')/I = O_2/I \setminus O'/I.$$

The reasoning is similar for both equalities, and therefore we restrict ourselves here to proving the first one; the second one can be proved by replacing O_2 by O'_1 . That $(O'_1 \setminus O')/I \supseteq O'_1/I \setminus O'/I$ follows from Lemma 36. To show that $(O'_1 \setminus O')/I \subseteq O'_1/I \setminus O'/I$, consider any \tilde{u} in $(O'_1 \setminus O')/I$, meaning that $\tilde{u} = [v_1]$ for some v in $O'_1 \setminus O'$. That implies already that $\tilde{u} \in O'_1/I$. Assume *ex absurdo* that $\tilde{u} \in O'/I = \tilde{O}'$, and consider any $v \in O'_1$ such that $[v] = \tilde{u}$, then also $v \in O_2$ and $[v] \in \tilde{O}'$, and therefore $v \in O'$. But this implies that $\tilde{u} \notin (O'_1 \setminus O')/I$, a contradiction. So we conclude that $(O'_1 \setminus O')/I = O'_1/I \setminus O'/I$ and that $(O_2 \setminus O')/I = O_2/I \setminus O'/I$, implying that

$$\tilde{O}'_1 \setminus \tilde{O}' = O'_1/I \setminus O'/I \subseteq (O_2/I \setminus O'/I) \setminus C/I(O_2/I \setminus O'/I) = (\tilde{O}_2 \setminus \tilde{O}') \setminus C/I(\tilde{O}_2 \setminus \tilde{O}'),$$

or equivalently, $(\tilde{O}'_1 \setminus \tilde{O}') \cap C/I(\tilde{O}_2 \setminus \tilde{O}') = \emptyset$. Since $\tilde{O}' \cap C/I(\tilde{O}_2 \setminus \tilde{O}') = \emptyset$, this also implies that $\tilde{O}'_1 \cap C/I(\tilde{O}_2 \setminus \tilde{O}') = \emptyset$. Since $\tilde{O}_2 \setminus \tilde{O}' = \tilde{O}_2 \setminus \tilde{O}$ and $\tilde{O}_1 \cap \tilde{O}_2 = \tilde{O}'_1$, this indeed leads to $\emptyset = \tilde{O}_1 \cap \tilde{O}_2 \cap C/I(\tilde{O}_2 \setminus \tilde{O}) = \tilde{O}_1 \cap C/I(\tilde{O}_2 \setminus \tilde{O})$.

For Axioms C₄a and C₄b, consider any \tilde{O}_1, \tilde{O}_2 in $\mathcal{Q}(\mathcal{V}/I)$ such that $\tilde{O}_1 \subseteq C/I(\tilde{O}_2)$, and consider any λ in $\mathbb{R}_{>0}$ and \tilde{u} in \mathcal{V}/I . Consider any option set O_2 in $\mathcal{Q}(\mathcal{V})$ such that $O_2/I = \tilde{O}_2$, let $O_1 := \{u \in O_2 : [u] \in \tilde{O}_1\} \in \mathcal{Q}(\mathcal{V})$, and consider any u in \mathcal{V} such that $[u] = \tilde{u}$. We first prove that then $O_1 \subseteq C(O_2)$. Indeed, consider any $v \in O_1$, meaning that $v \in O_2$ and $[v] \in \tilde{O}_1$, and therefore also $[v] \in C/I(\tilde{O}_2)$. Proposition 27 then guarantees that indeed $u \in C(O_2)$. We now infer from the coherence of C [Axioms C₄a and C₄b] that $\lambda O_1 \subseteq C(\lambda O_2)$ and $O_1 + \{u\} \subseteq C(O_2 + \{u\})$. Hence indeed $\lambda \tilde{O}_1 = \lambda O_1/I = (\lambda O_1)/I \subseteq C/I((\lambda O_2)/I) = C/I(\lambda O_2/I) = C/I(\lambda \tilde{O}_2)$ and $\tilde{O}_1 + \{\tilde{u}\} = O_1/I + \{[u]\} = (O_1 + \{u\})/I \subseteq C/I((O_2 + \{u\})/I) = C/I(O_2/I + \{[u]\}) = C/I(\tilde{O}_2 + \{\tilde{u}\})$, where the inclusions follow from Equation (14) and Proposition 27.

For the converse implication, assume that C/I is coherent. We show that C satisfies Axioms C₁–C₄b.

For Axiom C₁, consider any O in $\mathcal{Q}(\mathcal{V})$ and assume *ex absurdo* that $C(O) = \emptyset$. Then it follows from Proposition 27 that $C/I(O/I) = C(O)/I = \emptyset/I = \emptyset$, which contradicts the coherence of C/I [Axiom C₁].

For Axiom C₂, consider any u, v in \mathcal{V} such that $u < v$. Then $[u] < [v]$ by letting for example $w = 0$ in Proposition 23. The coherence of C/I [Axiom C₂] then guarantees that $\{[v]\} = C/I(\{[u], [v]\})$, implying that $\{v\} = C(\{u, v\})$, by Proposition 27.

For version (3) of Axiom C₃a, consider any option sets O, O_1, O_2 in $\mathcal{Q}(\mathcal{V})$ such that $O_1 \subseteq O_2 \subseteq O$ and $O_1 \cap C(O_2) = \emptyset$. Then Lemma 38 guarantees that $O_1/I \cap C/I(O_2/I) = \emptyset$. Because $O_2 \subseteq O$, we have by Equation (14) that $O_2/I \subseteq O/I$, whence $O_1/I \cap C/I(O/I) = \emptyset$ by coherence of C/I [version (3) of Axiom C₃a]. Using Lemma 38 again, we infer that indeed $O_1 \cap C(O) = \emptyset$.

For version (4) of Axiom C_{3b}, consider any option sets O, O_1, O_2 in $\mathcal{Q}(\mathcal{V})$ such that $O \subseteq O_1$ and $O_1 \cap C(O_2) = \emptyset$. We may assume without loss of generality that $O_1 \subseteq O_2$ [if not, replace O_1 by $O_1 \cap O_2$ and O by $O \cap O_2$]. Let

$$\bar{O}_1 := \{u \in O_2 : [u] \in O_1/I\} \supseteq O_1 \text{ and } \check{O} := \{u \in \bar{O}_1 : [u] \cap O_1 \subseteq O\},$$

then $\bar{O}_1/I = O_1/I$. Because C is compatible with I , we infer from Lemma 38 that the statements $O_1 \cap C(O_2) = \emptyset$, $\bar{O}_1 \cap C(O_2) = \emptyset$ and $\bar{O}_1/I \cap C/I(O_2/I) = \emptyset$ are equivalent, and therefore all hold. Observe that also $\check{O} \subseteq \bar{O}_1 \subseteq O_2$, implying that $\check{O}/I \subseteq \bar{O}_1/I \subseteq O_2/I$ by Equation (14). That $\check{O}/I \subseteq \bar{O}_1/I$ and $\bar{O}_1/I \cap C/I(O_2/I) = \emptyset$ implies, together with the coherence of C/I [version (4) of Axiom C_{3b}], that $\bar{O}_1/I \cap C/I(O_2/I \setminus \check{O}/I) = \emptyset$, and therefore also that $(\bar{O}_1/I \setminus \check{O}/I) \cap C/I(O_2/I \setminus \check{O}/I) = \emptyset$.

We now prove that $O_2/I \setminus \check{O}/I \subseteq (O_2 \setminus O)/I$. Consider any \tilde{u} in $O_2/I \setminus \check{O}/I$, and assume *ex absurdo* that $\tilde{u} \notin (O_2 \setminus O)/I$, or in other words, that $(\forall w \in O_2)(w \notin O \Rightarrow \tilde{u} \neq [w])$, or equivalently, $(\forall w \in O_2)(\tilde{u} = [w] \Rightarrow w \in O)$. This implies that $\tilde{u} \cap O_2 \subseteq O$, from which we infer on the one hand that $\tilde{u} \cap O_1 \subseteq O$ since $O_1 \subseteq O_2$ by assumption. On the other hand, we infer that $(\tilde{u} \cap O_2)/I \subseteq O/I$ by Equation (14), whence $\tilde{u} \in O/I$ because $\tilde{u} \in O_2/I$ by assumption, and therefore also $\tilde{u} \in O_1/I = \bar{O}_1/I$. Both observations together imply that $\tilde{u} \in \check{O}/I$, a contradiction with $\tilde{u} \in O_2/I \setminus \check{O}/I$. Coherence of C/I [version (3) of Axiom C_{3a}; with $\check{O}_1 := \bar{O}_1/I \setminus \check{O}/I$, $\check{O}_2 := O_2/I \setminus \check{O}/I$ and $\check{O} := (O_2 \setminus O)/I$] now implies that $(\bar{O}_1/I \setminus \check{O}/I) \cap C/I((O_2 \setminus O)/I) = \emptyset$.

We next prove that $(O_1 \setminus O)/I \subseteq \bar{O}_1/I \setminus \check{O}/I$. Consider any \tilde{u} in $(O_1 \setminus O)/I$, so $\tilde{u} = [v]$ for some v in $O_1 \setminus O$. This already implies that $\tilde{u} \in \bar{O}_1/I$. Assume *ex absurdo* that $\tilde{u} \notin \bar{O}_1/I \setminus \check{O}/I$, so $\tilde{u} \in \check{O}/I$, whence $\tilde{u} \cap O_1 \subseteq O$. But we know that $v \in O_1$ and $v \in \tilde{u}$, so $v \in O$, a contradiction. We conclude from all this that $(O_1 \setminus O)/I \cap C/I((O_2 \setminus O)/I) = \emptyset$, whence $O_1 \setminus O \cap C(O_2 \setminus O) = \emptyset$ by Lemma 38, and because $C(O_2 \setminus O) \subseteq O_2 \setminus O$, indeed that $O_1 \cap C(O_2 \setminus O) = \emptyset$.

For Axioms C_{4a} and C_{4b}, consider any O_1, O_2 in $\mathcal{Q}(\mathcal{V})$ such that $O_1 \subseteq C(O_2)$, and consider any λ in $\mathbb{R}_{>0}$ and \tilde{u} in \mathcal{V} . Then $O_1/I \subseteq C/I(O_2/I)$ by Equation (14) and Proposition 27, implying that $(\lambda O_1)/I = \lambda O_1/I \subseteq C/I(\lambda O_2/I) = C/I((\lambda O_2)/I)$ and $(O_1 + \{u\})/I = O_1/I + \{[u]\} \subseteq C/I(O_2/I + \{[u]\}) = C/I((O_2 + \{u\})/I)$, where the inclusions follow from the coherence of C/I [Axioms C_{4a} and C_{4b} respectively]. This then implies that indeed $\lambda O_1 \subseteq C(\lambda O_2)$ and $O_1 + \{u\} \subseteq C(O_2 + \{u\})$. \square

Proof of Proposition 29. Let us denote $C := \inf\{C_i : i \in \mathcal{I}\}$, then C is a coherent choice function by Proposition 8. We show that it is compatible with I . By assumption, and using Proposition 27, we have for all $i \in \mathcal{I}$ that

$$C_i(O) = \{u \in O : [u] \in C_i/I(O/I)\} \text{ for all } O \in \mathcal{Q}(\mathcal{V}).$$

Hence, for all $O \in \mathcal{Q}(\mathcal{V})$:

$$\begin{aligned} C(O) &= \bigcup_{i \in \mathcal{I}} C_i(O) = \bigcup_{i \in \mathcal{I}} \{u \in O : [u] \in C_i/I(O/I)\} \\ &= \{u \in O : (\exists i \in \mathcal{I})[u] \in C_i/I(O/I)\} \\ &= \{u \in O : [u] \in \bigcup_{i \in \mathcal{I}} C_i/I(O/I)\} \\ &= \{u \in O : [u] \in (\inf\{C_i/I : i \in \mathcal{I}\})(O/I)\}, \end{aligned}$$

and the stated result now follows from Propositions 8 and 27. \square

Lemma 40. *Given O_1 and O_2 in \mathcal{Q} such that $O_1 \subseteq O_2$, we have that*

$$O_1 \subseteq \{u \in O_2 : [u] \in O_1/I\}.$$

Proof. Consider any v in O_1 . Then $v \in O_2$ because $O_1 \subseteq O_2$, and $[v] \in O_1/I$ by the definition of the quotient space, whence $v \in \{u \in O_2 : [u] \in O_1/I\}$. \square

Proof of Proposition 30. We first prove that \approx_C is an equivalence relation. The reflexivity and symmetry are an immediate consequence of the definition (15). To prove transitivity, consider any u, v, w in \mathcal{V} and assume that $u \approx_C v$ and $v \approx_C w$. Consider any O in \mathcal{Q} and assume that $\{u, w\} \subseteq O$ and $u \in C(O)$. Then it suffices to prove that also $w \in C(O)$. Assume *ex absurdo* that $w \notin C(O)$, then we infer from Axiom C3a that also $w \notin C(O \cup \{v\})$. Since $v \approx_C w$, we infer from Equation (15) that then $v \notin C(O \cup \{v\})$, and similarly, since $u \approx_C v$, that $u \notin C(O \cup \{v\})$, so $\{u, v\} \subseteq (O \cup \{v\}) \setminus C(O \cup \{v\})$. Axiom C3b then tells us that $u \notin C(O)$, a contradiction.

To prove (ii), assume that $u \approx_C v$, and consider any $O \in \mathcal{Q}$ such that $\{u+w, v+w\} \subseteq O$ and $u+w \in C(O)$. Then it suffices to prove that also $v+w \in C(O)$. It follows from $u+w \in C(O)$ and Axiom C4b that $u \in C(O - \{w\})$. Since $u \approx_C v$, we infer from Equation (15) that then also $v \in C(O - \{w\})$, whence, again by Axiom C4b, indeed $v+w \in C(O)$.

To prove (i), assume that $u \approx_C v$. We first prove that then also $-u \approx_C -v$. Indeed, by applying (ii) with $w := -u - v$, we find that $-v \approx_C -u$. Now use the symmetry of \approx_C . Next, consider any $O \in \mathcal{Q}$ such that $\{\lambda u, \lambda v\} \subseteq O$ and $\lambda u \in C(O)$. Then it suffices to prove that also $\lambda v \in C(O)$. The proof is trivial if $\lambda = 0$. Because we have just proved that both $u \approx_C v$ and $-u \approx_C -v$, we may now assume without loss of generality that $\lambda > 0$. It follows from $\lambda u \in C(O)$ and Axiom C4a that $u \in C(\frac{1}{\lambda}O)$. Since $u \approx_C v$, we infer from Equation (15) that then also $v \in C(\frac{1}{\lambda}O)$, whence, again by Axiom C4a, indeed $\lambda v \in C(O)$.

We complete the proof by showing that I_C is a coherent set of indifferent options. To prove I₁, simply observe that $0 \approx_C 0$ by reflexivity of \approx_C . To prove I₂, it suffices to consider any $u \in \mathcal{V}_{>0}$, due to (i). It follows from Axiom C₂ that both $0 \notin C(\{0, u\})$ and $u \in C(\{0, u\})$, so we infer from Equation (15) that $0 \notin_C u$, whence indeed $u \in I_C$. To prove I₃, simply use (i). To prove I₄, simply use (ii) and the transitivity of \approx_C . \square

Proof of Proposition 31. For necessity, assume that C is compatible with I , and hence

$$(\forall O \in \mathcal{Q})C(O) = \{u \in O : [u] \in C(O)/I\}.$$

We need to prove that $I \subseteq I_C$, so consider any u in I and any $O \supseteq \{0, u\}$ in \mathcal{Q} such that $u \in C(O)$ [There always is such an O , for instance $O := \{0, u\}$, because $u \in I$ and therefore $[u] \in C/I(\{[u]\}) = C/I(\{0, u\}/I) = C(\{0, u\})/I$, by compatibility]. Then $[u] \in C(O)/I$ and because $[u] = [0]$, we find that $[0] \in C(O)/I$, so $0 \in C(O)$. Hence indeed $u \approx_C 0$, by Equation (15).

For sufficiency, assume that I is a linear subspace of I_C , so $I \subseteq I_C$, then we need to prove that C is compatible with I . In other words, if we consider any O in \mathcal{Q} and any $u \in O$ such that $[u] \in C(O)/I$, then we must prove that $u \in C(O)$. $[u] \in C(O)/I$ means that there is some $v \in C(O)$ such that $u - v \in I$, and therefore, by assumption, $u \approx_C v$. Since $\{u, v\} \subseteq O$ and $v \in C(O)$, we infer from Equation (15) that indeed $u \in C(O)$. \square

Proof of Lemma 32. Fix u and v in \mathcal{V} .

For the direct implication, assume that the left-hand side holds. We first prove that this implies that both $u - v \notin D$ and $v - u \notin D$. Indeed, consider the left-hand side for the particular choice $O := \{u, v\}$, leading to $\{0, v - u\} \cap D \neq \emptyset \Leftrightarrow \{0, u - v\} \cap D \neq \emptyset$, or equivalently, $v - u \in D \Leftrightarrow u - v \in D$, because $0 \notin D$ by Axiom D₁. So, if we had that $u - v \in D$ or $v - u \in D$, this would imply that both $u - v$ and $v - u$ would elements of D , and therefore also their sum $0 = u - v + v - u \in D$ by Axiom D₄. This contradicts Axiom D₁.

Next, consider any w in \mathcal{V} such that $w \notin \{u, v\}$, and consider the left-hand side for the particular choice $O := \{u, v, w\}$. Then $O \supseteq \{u, v\}$ and therefore $O - \{u\} \cap D \neq \emptyset \Leftrightarrow O - \{v\} \cap D \neq \emptyset$, which can be rewritten as $\{0, v-u, w-u\} \cap D \neq \emptyset \Leftrightarrow \{u-v, 0, w-v\} \cap D \neq \emptyset$. But $0 \notin D$ by Axiom D_1 , whence $\{v-u, w-u\} \cap D \neq \emptyset \Leftrightarrow \{u-v, w-v\} \cap D \neq \emptyset$. But, since we have seen above that neither $u-v \in D$ nor $v-u \in D$, this can in turn be rewritten as $w-u \notin D \Leftrightarrow w-v \notin D$. This tells us that $w \in D + \{u\} \Leftrightarrow w \in D + \{v\}$ for all w in \mathcal{V} , and therefore indeed $D + \{u\} = D + \{v\}$.

For the converse implication, assume that $D + \{u\} = D + \{v\}$. This immediately allows us to infer that $u-v \notin D$ and $v-u \notin D$. Consider any $O \supseteq \{u, v\}$. If $O = \{u, v\}$ then $O - \{u\} \cap D = \emptyset$ and $O - \{v\} \cap D = \emptyset$. Let $n := |O| - 2$ and assume therefore that $n \geq 1$. Label the elements of $O \setminus \{u, v\}$ as w_1, \dots, w_n , without loss of generality. Then $O = \{u, v, w_1, \dots, w_n\}$. Infer the following equivalences:

$$\begin{aligned} O - \{u\} \cap D \neq \emptyset &\Leftrightarrow \{0, v-u, w_1-u, \dots, w_n-u\} \cap D \neq \emptyset \\ &\Leftrightarrow \{w_1-u, \dots, w_n-u\} \cap D \neq \emptyset \\ &\Leftrightarrow (\exists i \in \{1, \dots, n\}) w_i - u \in D \\ &\Leftrightarrow (\exists i \in \{1, \dots, n\}) w_i \in D + \{u\} = D + \{v\} \\ &\Leftrightarrow (\exists i \in \{1, \dots, n\}) w_i - v \in D \Leftrightarrow O - \{v\} \cap D \neq \emptyset, \end{aligned}$$

where the second equivalence follows from the fact that $0 \notin D$ [see Axiom D_1] and our earlier observation that $v-u \notin D$. \square

Proof of Proposition 33. We begin with the first statement. Consider any coherent choice function C that is compatible with I . We must prove that $D_C + I \subseteq D_C$. Observe that for any w in \mathcal{V} :

$$\begin{aligned} w \in D_C &\Leftrightarrow 0 \notin C(\{0, w\}) \Leftrightarrow [0] \notin C/I(\{[0], [w]\}) \\ &\Leftrightarrow [w] \in D_{C/I}, \end{aligned} \tag{17}$$

where the first equivalence follows from Proposition 12 and the second from the compatibility of C with I . So, consider any v in D_C and any u in I , then $[0] \notin C/I(\{[0], [v]\})$ and $[u] = [0] = I$, implying that $[v+u] = [v] + [u] = [v] + I = [v] + [0] = [v+0] = [v]$, and therefore $[0] \notin C/I(\{[0], [v+u]\})$, implying that indeed $v+u \in D_C$. The last statement follows directly from Equation (17).

We turn now towards the second statement. Consider any coherent set of indifferent options I and any coherent set of desirable options D such that $D + I \subseteq D$. We must prove that $C_D(O) = \{u \in O : [u] \in C_D/I(O/I)\}$ for all O in $\mathcal{Q}(\mathcal{V})$. Due to Proposition 13, we know that $C_D(O) = \{u \in O : (\forall v \in O) v-u \notin D\}$ for all O in $\mathcal{Q}(\mathcal{V})$, whence, by Proposition 24,

$$\begin{aligned} C_D(O) &= \{u \in O : (\forall v \in O) [v] - [u] \notin \{[w] : w \in D\}\} \\ &= \{u \in O : (\forall v \in O) [v] - [u] \notin D/I\} \\ &= \{u \in O : (\forall [v] \in O/I) [v] - [u] \notin D/I\} \\ &= \{u \in O : [u] \in C_{D/I}(O/I)\}, \end{aligned}$$

for all O in $\mathcal{Q}(\mathcal{V})$, because $C_{D/I}(O/I) = \{\tilde{u} \in O/I : (\forall \tilde{v} \in O/I) \tilde{v} - \tilde{u} \notin D/I\}$ for all O/I in $\mathcal{Q}(\mathcal{V}/I)$. The fact that $C_D/I = C_{D/I}$ now follows from Proposition 27. \square

Proof of Proposition 34. That the transformation $\text{inv}_{\mathcal{P}}$ is linear, is immediate from its definition in Equation (16).

To prove (i), observe that $\text{inv}_{\mathcal{P}} \circ \pi^t = \frac{1}{|\mathcal{P}|} \sum_{\varpi \in \mathcal{P}} \varpi^t \circ \pi^t = \frac{1}{|\mathcal{P}|} \sum_{\varpi \in \mathcal{P}} (\pi \circ \varpi)^t = \text{inv}_{\mathcal{P}}$, where the last equality holds because \mathcal{P} is a group. For the second identity, observe that $\pi^t \circ \text{inv}_{\mathcal{P}} = \frac{1}{|\mathcal{P}|} \sum_{\varpi \in \mathcal{P}} \pi^t \circ \varpi^t = \frac{1}{|\mathcal{P}|} \sum_{\varpi \in \mathcal{P}} (\varpi \circ \pi)^t = \text{inv}_{\mathcal{P}}$, where the first equality follows from the linear character of π^t and the last equality holds because \mathcal{P} is a group.

To prove (ii), observe that

$$\text{inv}_{\mathcal{P}} \circ \text{inv}_{\mathcal{P}} = \frac{1}{|\mathcal{P}|} \sum_{\pi \in \mathcal{P}} \text{inv}_{\mathcal{P}} \circ \pi^t = \frac{1}{|\mathcal{P}|} \sum_{\pi \in \mathcal{P}} \text{inv}_{\mathcal{P}} = \text{inv}_{\mathcal{P}},$$

where the first equality is due to the linearity of $\text{inv}_{\mathcal{P}}$ and the second due to (i).

To prove (iii), consider any u in $\text{kern}(\text{inv}_{\mathcal{P}})$. Then $\text{inv}_{\mathcal{P}} u = 0$ and therefore $u = u - \text{inv}_{\mathcal{P}} u = \frac{1}{|\mathcal{P}|} \sum_{\pi \in \mathcal{P}} (u - \pi^t u)$ is an element of $I_{\mathcal{P}}$. Conversely, consider any u in $I_{\mathcal{P}}$, then $u = \sum_{k=1}^n \lambda_k (v_k - \pi_k^t v_k)$ for some n in \mathbb{N} , λ_k in \mathbb{R} , v_k in \mathcal{V} and π_k in \mathcal{P} . But then $\text{inv}_{\mathcal{P}} u = \sum_{k=1}^n \lambda_k (\text{inv}_{\mathcal{P}} v_k - \text{inv}_{\mathcal{P}} (\pi_k^t v_k)) = 0$, where the first equality is due to the linearity of $\text{inv}_{\mathcal{P}}$ and the last due to (i). Hence indeed $u \in \text{kern}(\text{inv}_{\mathcal{P}})$.

To prove (iv), consider any u in $\text{rng}(\text{inv}_{\mathcal{P}})$. Then $u = \text{inv}_{\mathcal{P}} v$ for some v in \mathcal{V} , and therefore $\pi^t u = \pi^t (\text{inv}_{\mathcal{P}} v) = (\pi^t \circ \text{inv}_{\mathcal{P}}) v = \text{inv}_{\mathcal{P}} v = u$ for all π in \mathcal{P} , where the third equality follows from (i). Hence indeed $u \in \mathcal{V}_{\mathcal{P}}$. Conversely, consider any u in $\mathcal{V}_{\mathcal{P}}$. Then $\pi^t u = u$ for all π in \mathcal{P} , and therefore $u = \text{inv}_{\mathcal{P}} u$, whence indeed $u \in \text{rng}(\text{inv}_{\mathcal{P}})$.

For the last statement, simply observe that $v \in u/I_{\mathcal{P}} \Leftrightarrow v - u \in I_{\mathcal{P}} \Leftrightarrow \text{inv}_{\mathcal{P}}(v - u) = 0 \Leftrightarrow \text{inv}_{\mathcal{P}} v = \text{inv}_{\mathcal{P}} u$, where the second equivalence follows from (iii) and the last from the linearity of $\text{inv}_{\mathcal{P}}$. \square

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